

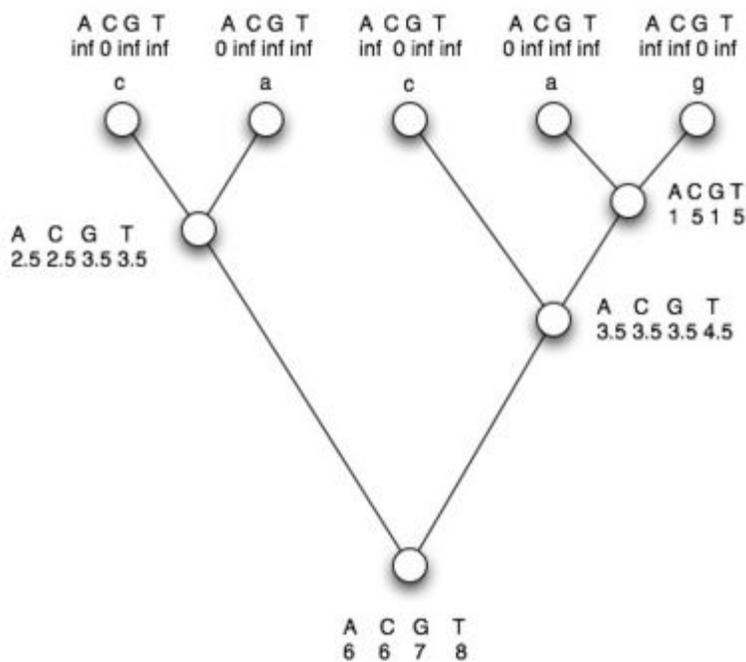
20.181 Lecture 5

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Downpass (cont.)

The tree we were working on last time:
Now we're going to try to code up the downpass.



1. So now we want to know: what is the sequence at this ancestor node, that we can't observe? We'll try all four possibilities, and calculate the penalty associated with each of those possibilities are.
2. Remember that the good solution with the fibonacci series was to pass all the information back each time so that we don't have to repeat calculations!
3. We left off at the root node, where we said there were 64 possibilities.
4. So let's go through this by hand and try to write pseudo code for a function that will do the same thing.

5. What we're going to do, which might seem odd at first, is to pass the node's identity (nodeSeq) (A,C,G,or T) into the function as if we know it. Our function will look at this one possibility and return one column of the cost vector. We will pass to this function *each* possibility, one by one.

Sankoff Downpass Algorithm

(See lecture 8 notes for a more efficient way of implementing Sankoff downpass)

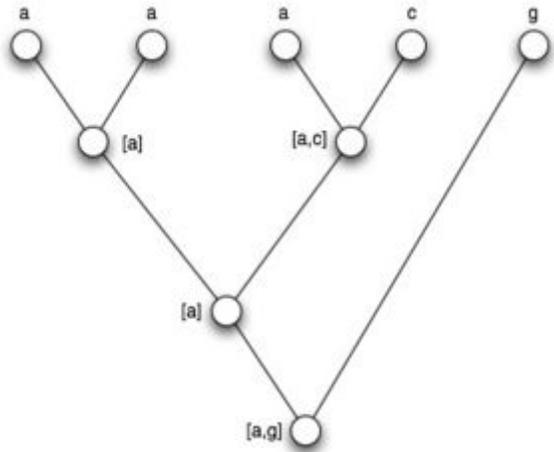
	A	C	G	T
A	0	2.5	1	2.5
C	2.5	0	2.5	1
G	1	2.5	0	2.5
T	2.5	1	2.5	0

```
def sankoff(tree,nodeSeq):
    min = inf #initialization
    if tree is a leaf: #stop condition
        if tree['data'] == nodeSeq:
            return 0
        return inf
    for leftSeq in [A,C,G,T]: #now we're at a node with
        two children
        for rightSeq in [A,C,G,T]: #we have to try all 16
            possibilities here
            sum= cost(nodeSeq,leftSeq,rightSeq) #cost is some
            function that looks up the cost in the table
            sum += sankoff(tree['left'],leftSeq) #we have to
            remember to pass the cost along!
            sum += sankoff(tree['right'],rightSeq)
            if (sum < min):
                min = sum
    return min
```

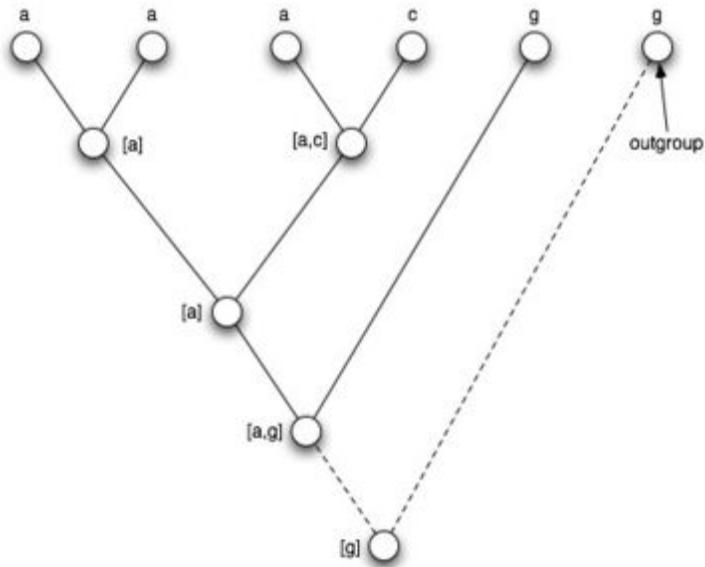
```
def sankoffDownpass(tree): #we don't know what the sequence
    is at the root

    for seq in [A,C,G,T]:
        if sankoff(tree,seq) < min:
            min = sankoff(tree,seq)
            minSeq = seq
    return min
```

Downpass: can we stop there?



- So far we're only passing the best score obtained at each subtree down the tree to the root.
- The downPass algorithm, which you'll implement in HW4, finds the most parsimonious sequence at the root node.
- Why does it find the most parsimonious sequence **only** for the root ?



- - Think about what the difference is between the root and the other internal nodes:

the root doesn't have any parents -- only children.

- So far, all our information flows down tree, and we use it to make inference about older and older internal nodes. That's why the inference at the root the best

possible guess we can make; it has **all** the information from **everywhere** on the tree. But the inferences we've written at the internal nodes are NOT YET correct- these internal nodes only have *part* of the information on the tree.

- What if we added an outgroup, such that what *was* the root before is now an internal node
 - would this change our guess about nucleotides that are likely at the old root?

yes! because we would have additional information that had not been taken into account in our prior best guess

Fitch's upPass

We'll pass the information back up from the root node to the other internal nodes in order to decide what their most likely sequences were!

- computer science jargon: three ways to traverse trees
 1. pre-order: do some operation, then call left subtree, then call right subtree

like typing into a calculator: + 5 3

2. in-order: L, op, R

like typing into a calculator: 5 + 3

3. post-order: L, R, op

like typing into a calculator: 5 3 +