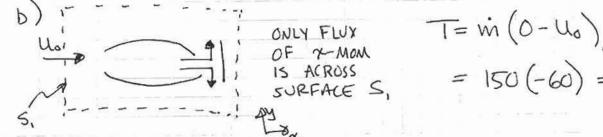


WAITZ (10/1)



$$T = m(0 - u_0)$$
  
= 150(-60) =  $9 \text{ kN}$ 

C) WHEN AIRPLANE COMES TO REST, THERE IS NO NET FORCE IN THE X-DIR. YOU CAN SEE THIS BY DRAWING THE CONTROL VOUNCE LARGE RELATIVE THE ENGINE

INTEGRAL MOMENTUM THEOREM IN FRAME ATTACHED TO VEHICLE:

# UNIFIED PROPULSION PZ SOLUTIONS

WAITZ

USE INTEGRAL MUMENTUM THEOREM IN VEHICLE REFERENCE FRAME.

VELOCITY OF REFERENCE FRAME = V (VELOCITY OF VEHICLE).

ALL OTHER VELOCITIES ARE RELATIVE TO V AND ARE LABELED.

The continuent of th

 $\sum_{i=1}^{n} F_{i}$ 

and the second of the second o

= dy(sdte

SUM OF EXTERNAL FORLES = O SINCE PROBLEM SAYS TO NEGLECT DRAG

ACCELSRATION RECATIVE TO INSERTIAL REF. FRAME

= Malt

- (d(pux) dVoe + Sux(pu). rids Voe s

NET

FLUX OF

MON.

BOUNDAMES

ACROSS

CHANGE IN FROM.
OF MASS WITHIN C.V.
RELATIVE TO C.V.
REF. FRAME

= 0 SINCE

MASS OF VEHICLE NOT CHANGING AND CAN NEGLECT ANY UNSTEADY CHANGES

IN MUMENTUM OF WATER PER-PROBLEM

STATEMENT

C.V. MOVING WITH

AIRPLANE AT

V(t)

 $0 = M \frac{dV}{dt} + \int_{S} u_{x} (g\vec{u}) \cdot \vec{n} dS$ 

SOF 2)

Flow in

Flow out

OPPOSITE
TO NORMAL

X-DIRECTION

NOTE THE SIGNS ARE BOTH NEGATIVE BUT FOR DIFFERENT REASONS!

EQNS OF MOTION: 
$$\frac{dV}{dt} = \frac{2Ajfi}{M} (Vj - V)^2 \qquad Vj > V$$

$$\frac{dV}{dt} = 0$$

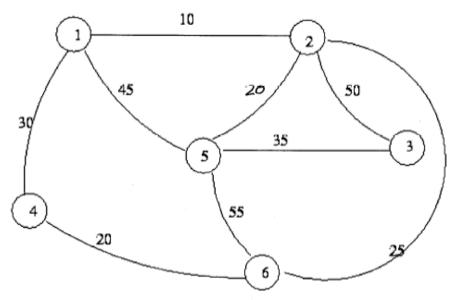
$$Vj < V$$

c) 
$$F = -2g_i A_i (V_i - V)^2$$
  
 $T = -F = 2g_i A_i (V_i - V)^2$ 

#### Home Work 9

The problems in this problem set cover lectures C7, C8, C9 and C10

1. What is the Minimum Spanning Tree of the graph shown below using both Prim's and Kruskal's algorithm. Show all the steps in the computation of the MST (not just the final MST).

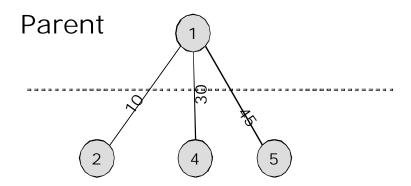


### **Prim's Algorithm**

Step 1.

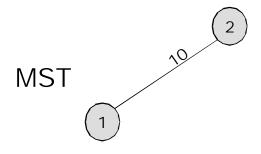
**MST** 

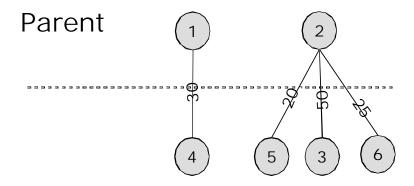




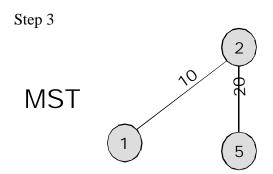
Fringe Set

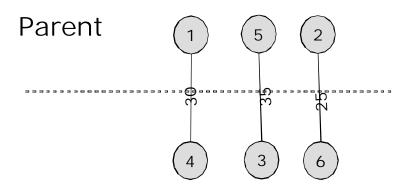
Step 2



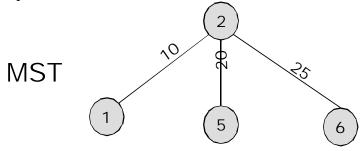


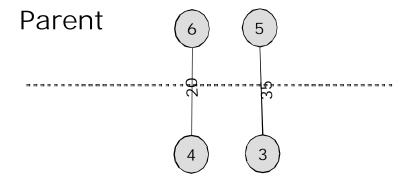
## Fringe Set



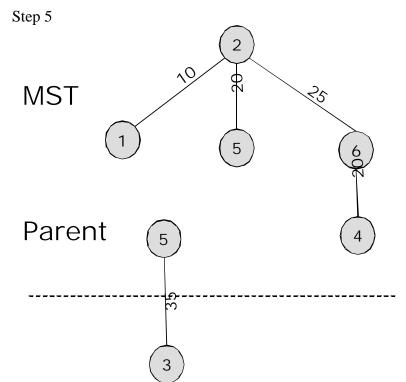


Fringe Set



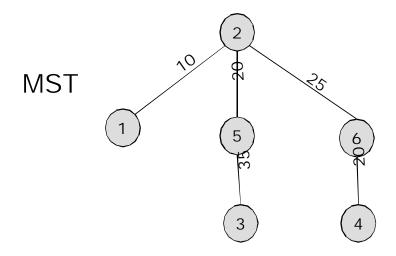


Fringe Set



Fringe Set

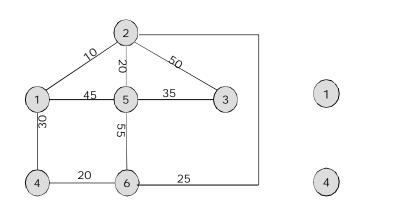
## Minimum Spanning Tree

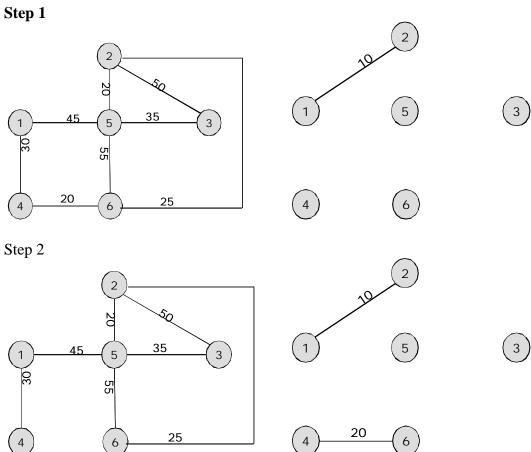


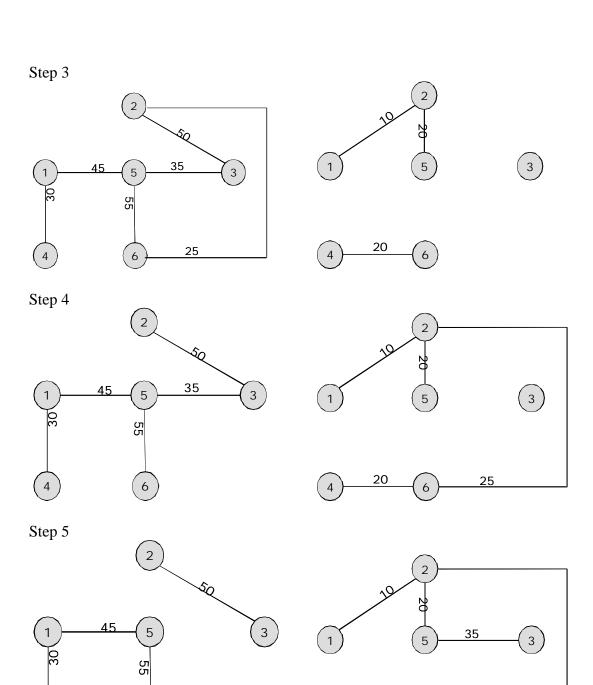
Weight of the MST = 10 + 20 + 25 + 20 + 35 = 110

## Kruskals Algorithm

## Initialization







Weight of the MST = 10 + 20 + 20 + 25 + 35 = 110

2. Compute the computation complexity of the bubble sort algorithm. Show all the steps in the computation based on the algorithm.

#### Algorithm

Procedure Bubble\_Sort(Input\_Output\_Array)

end Bubble\_Sort;

$$O(N(N-1)) = O(N^2)$$

- 3. What are the best case and worst case computation complexity of:
- a. Inserting a node into an unsorted singly linked list

Inserting into an unsorted singly linked list is carried out using the add\_to\_front operation.

Both the best and worst case execution time is O(1).

b. Inserting a node into a sorted singly linked list

In the case of a sorted linked list, the list has to be traversed to find the right position. The list traversal takes O(n) in the worst case.

Best case execution time is O(1) if the element being inserted is the smallest element in the list (list in ascending order)

Worst case execution time is O(n) if the element being inserted is the largest element in the list (list in ascending order)

4.

- a. Design an Ada95 Package to:
  - i. Read in N integers from an input file into an array. (N is user defined <=50)
  - ii. Sort the array in ascending order

type my\_array is array (1 .. 50) of integer;

iii. Perform binary search on the array.

The package is designed as follows: Data Structures:

```
Subprograms:
```

```
-- procedure to create an array with <=50 elements
-- assumes that input can be found in input_file.txt
-- accepts the array
-- returns array with elements read from file and the number of elements read
procedure Create (
   Num_Array : in out My_Array;
   Size : out Integer );
procedure Merge (
   Input_Array : in out My_Array;
   Lb_1
            : in Integer;
   Ub_1
            : in Integer;
   Lb 2
             : in Integer;
   Ub 2
             : in Integer );
-- procedure to sort the array in ascending order
-- accepts the array and the lowerbound(lb) and upperbound(ub)
-- returns the sorted array
procedure Merge_Sort (
   Num_Array: in out My_Array;
   Lb
          : in Integer;
   Ub
          : in Integer );
--function to perform binary search on the array
-- accepts the array, lb, ub and the element being searched for
-- returns the index of the element if it is found
-- returns -1 if the element is not found
function Binary_Search (
   Num_Array : My_Array;
   Lb
           : Integer;
   Ub
            : Integer;
   Looking_For : Integer )
 return Integer;
```

#### Algorithms:

#### Create

Pre-conditions: An array of 50 elements (num\_array)

Post-Condition: Array with a maximum of 50 elements loaded from the file, size of the

array.

#### Algorithm

- 1. Open the file input\_file.txt
- 2. Initialize the counter to 0
- 3. While not End\_Of\_File (input\_file)
  - a. Increment the counter
  - b. Read an element from the file into num\_array(Counter)
- 4. Return num\_array and Counter

#### Merge\_Sort

Merge Sort is a *sort* algorithm that splits the items to be sorted into two groups, *recursively* sorts each group, and *merges* them into a final, sorted sequence.

Pre-Conditions: An array of 1 or more elements

Post-Condition : A sorted array

#### Algorithm

- 1. Check if LB < UB
  - a. Call merge sort with
    - 1. Input Array
    - LB = LB
    - 3. UB = (UB + LB)/2
  - b. Call merge sort with
    - 1. Input Array
    - 2. LB = (UB + LB)/2 + 1
    - 3. UB = UB
  - c. Call merge with
    - 1. Merge with parameters
      - Array to be sorted
      - Lower bound
      - (Lower bound + Upper bound) / 2
      - (Lower bound + Upper bound) / 2 +1
      - Upper bound

The merge sort procedure will recursively call itself until it has only single element arrays. These are inherently sorted. It will then merge the successively larger elements until the whole sorted array is produced.

#### Merge

```
Pre-Conditions: An array of 1 or more elements, Legal values of Lower_Bound_1, Upper_Bound_1, Lower_Bound_2 and Upper_Bound_2,
```

Post-Condition: A merged array that is sorted

#### Algorithm

- 1. Check if
  - a. Lower\_Bound\_1 < Upper\_Bound\_1
  - b. Lower\_Bound\_2 < Upper\_Bound\_2
  - c. Lower\_Bound\_1 < Lower\_Bound\_2
  - d. Upper\_Bound\_1 < Upper\_Bound\_2
- 2. If any of the above conditions are violated,
  - a. Display Error
  - b. Stop Executing the program
- 3. Set
  - a. Index 1 := Lower Bound 1;
  - b. Index\_2 := Lower\_Bound\_2;
  - c. Index := Lower Bound 1;
  - d. Temp\_Array := Input\_Array
- 4. While (Index\_1 <= Upper\_Bound\_1) and (Index\_2 < = Upper\_Bound\_2)
  - a. If (Input Array(Index 1) < Input Array (Index 2)
    - i. Temp\_Array(Index) := Input\_Array(Index\_1)
    - ii.  $Index_1 := Index_1 + 1$ ;
    - iii. Index := Index + 1;
  - b. Else
    - i. Temp\_Array(Index) := Input\_Array(Index\_2)
    - ii. Index 2 := Index 2 + 1:
    - iii. Index := Index + 1;
- 5. While (Index\_1 <= Upper\_Bound\_1)
  - a. Temp\_Array(Index) := Input\_Array(Index\_1)
  - b.  $Index_1 := Index_1 + 1$ ;
  - c. Index := Index + 1;
- 6. While (Index\_2 <= Upper\_Bound\_2)
  - a. Temp\_Array(Index) := Input\_Array(Index\_2)
  - b. Index 2 := Index 2 + 1;

```
c. Index := Index + 1;
```

7. Input\_Array := Temp\_Array

#### **Binary Search**

Pre-Conditions: Array to be searched

Item that you are searching for

Post-condition: Index location of the item you are searching for

Return -1 if the number is not found.

Algorithm

1. Set Return\_Index to -1;

- 2. Set Current\_Index to (Upper\_Bound Lower\_Bound + 1) /2.
- 3. Loop
- i. if the lower\_bound > upper\_boundExit:
- ii. if ( Input\_Array(Current\_Index) = Number\_to\_Search) then
   Return\_Index = Current\_Index)
   Exit;
- iii. if ( Input\_Array(Current\_Index) > Number\_to\_Search) then
   Lower\_Bound = Current\_Index +1
   else
   Upper\_Bound = Current\_Index 1
- 4. Return Return Index
- b. Write a program to test your package that will
  - Prompt the user for a number to search for.
  - If the number is found using the binary search algorithm
    - Display the location (index)
    - Display the number
  - If the number is not found using the binary search algorithm
    - Display "Number not in array to the user"

```
My_Test_Array: My_Array;
 Size
           : Integer;
 Location : Integer;
 Find: Integer;
begin
  -- create the array
 Create(My_Test_Array,Size);
  -- get number to search for from user
 Ada.Text_Io.Put("Please Enter the Number to Search For");
 Ada.Integer_Text_Io.Get(Find);
 -- display unsorted array to the user
 for I in 1.. Size loop
   Ada.Text_Io.Put_Line(Integer'Image(My_Test_Array(I)));
 end loop;
 Ada.Text_Io.New_Line;
 -- sort the array using the merge sort algorithm
 Merge_Sort(My_Test_Array,1,Size);
 -- display the sorted algorithm
 for I in 1..Size loop
   Ada.Text_Io.Put_Line(Integer'Image(My_Test_Array(I)));
 end loop;
 -- perform a binary search on the array
 Location:= Binary_Search(My_Test_Array, 1, Size, Find);
 if Location /= -1 then
   Ada.Text_Io.Put("Found number at");
   Ada.Text_Io.Put_Line(Integer'Image(Location));
   Ada.Text_Io.Put_Line(Integer'Image(My_Test_Array(Location)));
   Ada.Text_Io.Put_Line("Number Not Found in Array");
 end if;
end Test_Hw_9;
```

- 5. Implement the merge sort algorithm as an Ada95 program. Your program should
  - Read in N integers from an input file. (N is user defined <=50)
  - Sort using your merge sort implementation.
  - Display the sorted and unsorted inputs to the user

Solved in problem 4.

#### **Problem S8 Solution**

1. The convolution is given by

$$y(t) = g(t) * u(t) = \int_{-\infty}^{\infty} g(t - \tau)u(\tau) d\tau$$
 (1)

Note that  $u(\tau)$  is nonzero only for  $-3 \le \tau \le 0$ , and  $g(t-\tau)$  is nonzero only for  $0 \le t - \tau \le 3$ , that is, for  $-3 + t \le \tau \le t$ . So there are four distinct regimes:

- (a) t < -3
- (b)  $-3 \le t \le 0$
- (c) 0 < t < 3
- (d) t > 3

For cases (a) and (d), there is no overlap between  $g(t - \tau)$  and  $u(\tau)$ , so y(t) = 0. For case (b), the overlap is for  $-3 \le \tau \le t$ . So

$$y(t) = \int_{-\infty}^{\infty} g(t-\tau)u(\tau) d\tau$$
$$= \int_{-3}^{t} \sin(-2\pi(t-\tau))\sin(2\pi\tau) d\tau$$

At this point, we have to do a little trig:

$$\sin(-2\pi(t-\tau))\sin(2\pi\tau) = \sin(2\pi(\tau-t))\sin(2\pi\tau) 
= [\sin(2\pi\tau)\cos(2\pi t) - \cos(2\pi\tau)\sin(2\pi t)]\sin(2\pi\tau) 
= \cos(2\pi t)\sin^2(2\pi\tau) - \sin(2\pi t)\cos(2\pi\tau)\sin(2\pi\tau) 
= \cos(2\pi t)\frac{1-\cos(4\pi\tau)}{2} - \sin(2\pi t)\frac{\sin(4\pi\tau)}{2}$$

So the integral is given by

$$y(t) = \int_{-3}^{t} \frac{\cos(2\pi t)}{2} d\tau - \int_{-3}^{t} \frac{\cos(2\pi t)}{2} \cos(4\pi \tau) d\tau - \int_{-3}^{t} \frac{\sin(2\pi t)}{2} \sin(4\pi \tau) d\tau$$

$$= \frac{\cos(2\pi t)}{2} (t+3) - \frac{\cos(2\pi t)}{8\pi} \sin(4\pi \tau) \Big|_{\tau=-3}^{t} + \frac{\sin(2\pi t)}{8\pi} \cos(4\pi \tau) \Big|_{\tau=-3}^{t}$$

$$= \frac{\cos(2\pi t)}{2} (t+3) - \frac{\cos(2\pi t)}{8\pi} \sin(4\pi t) + \frac{\sin(2\pi t)}{8\pi} [\cos(4\pi t) - 1]$$

(As often happens with problems involving trig functions, there are other equivalent expressions.)

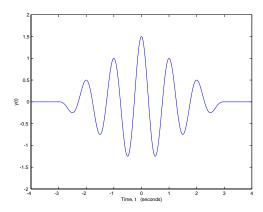
For case (c), the region of integration is  $-3 + t \le \tau \le 0$ . So

$$y(t) = \int_{-3+t}^{0} \frac{\cos(2\pi t)}{2} d\tau - \int_{-3+t}^{0} \frac{\cos(2\pi t)}{2} \cos(4\pi \tau) d\tau - \int_{-3+t}^{0} \frac{\sin(2\pi t)}{2} \sin(4\pi \tau) d\tau$$

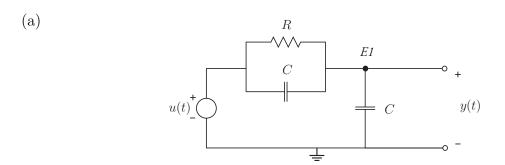
$$= \frac{\cos(2\pi t)}{2} (3-t) - \frac{\cos(2\pi t)}{8\pi} \sin(4\pi \tau) \Big|_{\tau=-3+t}^{0} + \frac{\sin(2\pi t)}{8\pi} \cos(4\pi \tau) \Big|_{\tau=-3+t}^{0}$$

$$= \frac{\cos(2\pi t)}{2} (3-t) + \frac{\cos(2\pi t)}{8\pi} \sin(4\pi t) - \frac{\sin(2\pi t)}{8\pi} [\cos(4\pi t) - 1]$$

2. y(t) is plotted below.



- 3. The maximum value of y(t-T) occurs at time T. So I would use this center peak to identify the delay time T.
- 4. The adjacent peaks are nearly as tall as the center peak, so if noise were added to the signal, the tallest peak might not be the center peak, so we might use the wrong peak to determine the delay time.
- 5. The chirp signal of Problem S6 produces an ambiguity function with only one prominent peak. Therefore, the addition of noise should not make it difficult to accurately determine the delay time.



We can use impedance methods to solve for Y(s) in terms of U(s). Label ground and  $E_1$  as shown. Then KCL at  $E_1$  yields

$$Cs(E_1 - 0) + \left(Cs + \frac{1}{R}\right)(E_1 - U) = 0$$
 (1)

Simplifying, we have

$$\left(2Cs + \frac{1}{R}\right)E_1 = \left(Cs + \frac{1}{R}\right)U$$

Since we are finding the step response,

$$U(s) = \frac{1}{s}, \quad \operatorname{Re}[s] > 0$$

Plugging in numbers, we have

$$(0.5s + 0.5)E_1(s) = (0.25s + 0.5)\frac{1}{s}$$

Solving for  $E_1$ , we have

$$E_1(s) = \frac{0.25s + 0.5}{(0.5s + 0.5)s} = \frac{0.5s + 1}{s(s+1)}$$

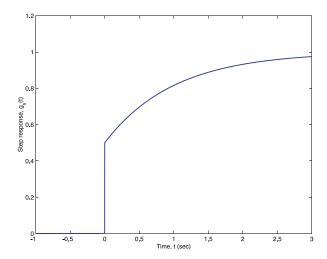
The region of convergence must be Re[s] > 0, since the step response is causal, and the pole at s = 0 is the rightmost pole. Using partial fraction expansions,

$$E_1(s) = \frac{1}{s} - \frac{0.5}{s+1}$$

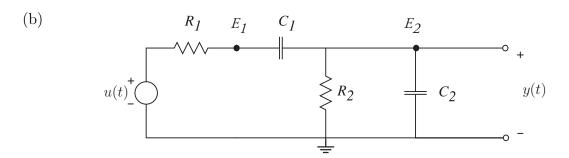
Therefore,  $g_s(t) = y(t) = e_1(t)$  is the inverse transform of  $E_1(t)$ , so

$$y(t) = \left(1 - 0.5e^{-t}\right)\sigma(t)$$

The step response is plotted below:



Normal differential equation methods are difficult to apply, because we cannot apply the normal initial condition that  $e_1(0) = 0$ . This is because the chain of capacitors running from the voltage source to ground causes there to be an impulse of current at time t = 0, and the voltages across the capacitors change instantaneously at t = 0. It is possible to use differential equation methods, we just have to be more careful about the initial conditions. However, Laplace methods are easier.



Again, use impedance methods, using the node labelling above. Then the node equations are

$$(C_1s + G_1)E_1 - C_1sE_2 = G_1U$$

$$-C_1sE_1 + [(C_1 + C_2)s + G_2]E_2 = 0$$
(2)

where G = 1/R. We can use Cramer's rule to solve for  $E_2$ :

$$E_{2}(s) = \frac{\begin{vmatrix} C_{1}s + G_{1} & G_{1}U(s) \\ -C_{1}s & 0 \end{vmatrix}}{\begin{vmatrix} C_{1}s + G_{1} & -C_{1}s \\ -C_{1}s & (C_{1} + C_{2})s + G_{2} \end{vmatrix}}$$

$$= \frac{G_{1}C_{1}s}{C_{1}C_{2}s^{2} + (G_{1}C_{1} + G_{1}C_{2} + G_{2}C_{1})s + G_{1}G_{2}}U(s)$$
(3)

Since we are finding the step response,

$$U(s) = \frac{1}{s}, \quad \text{Re}[s] > 0$$

Plugging in numbers, we have

$$Y(s) = E_2(s) = \frac{0.1s}{0.06s^2 + 0.35s + 0.25} \frac{1}{s} = \frac{5/3}{s^2 + 5.833\overline{3}s + 4.166\overline{6}}$$
(4)

In order to find y(t), we must expand Y(s) in a partial fraction expansion. To do so, we must factor the denominator, using either numerical techniques or the quadratic formula. The result is

$$s^{2} + 5.833\bar{3}s + 4.166\bar{6} = (s+5)(s+0.833\bar{3})$$
 (5)

We can use the coverup method to factor Y(s), so that

$$Y(s) = \frac{5/3}{(s+5)(s+0.833\overline{3})} = \frac{-0.4}{s+5} + \frac{0.4}{s+0.833\overline{3}}$$
 (6)

The region of convergence must be  $\text{Re}[s] > -0.833\overline{3}$ , since the step response is causal, and the r.o.c. is to the right of the right-most pole. Therefore, the step response is given by the inverse transform of Y(s), so that

$$g_s(t) = \left(-0.4e^{-5t} + 0.4e^{-0.833\bar{3}t}\right)\sigma(t) \tag{7}$$

The step response is plotted below:

