# Introduction to Computers and Programming 

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## Numeric Values

- Storing the value of $25_{10}$ using ASCII: 0011001000110101
- Binary notation: $00000000000011001_{2}$



## Finding Binary Representation of Large Values

1. Divide the value by 2 and record the remainder
2. As long as the quotient obtained is not 0 , continue to divide the newest quotient by 2 and record the remainder
3. Now that a quotient of 0 has been obtained, the binary representation of the original value consists of the remainders listed from right to left in the order they were recorded


The Binary System

- Decimal: Position represents a power of 10

$$
-5382_{10}=5 \times 10^{3}+3 \times 10^{2}+8 \times 10^{1}+2 \times 10^{0}
$$

- Binary: Position represents a power of 2
$-1011_{2}=1 \times 2^{3}+0 \times 2^{2}+1 \times 2^{1}+1 \times 2^{0}$
$=8+0+2+1=11_{10}$
- Binary addition

$$
\begin{array}{rrrr}
0 & 1 & 0 & 1 \\
+0 & +0 & +1 & +1 \\
\hline 0 & +1 & 1 & +1
\end{array}
$$

- A byte
bit number

bit value

$$
\begin{array}{llllllll}
128 & 64 & 32 & 16 & 8 & 4 & 2 & 1
\end{array}
$$

## Representing Negative Numbers

- Using one byte, any suggestions for representing negative numbers?
- E.g., $00010011_{2}=19_{10}$. How to represent $19_{10}$ in binary?
- Reserve 1 bit (\#7) for sign, 7 bits for number (sign magnitude)
$-\mathbf{0 0 0 0 0 0 0 0 1}_{2}=1_{10} \quad \mathbf{1 0 0 0 0 0 0 1}_{2}=-1_{10}$
$-00010011_{2}=19_{10} \quad 10010011_{2}=-19_{10}$


## Representing Negative Numbers

- One's complement - invert each bit
$-00010011_{2}=19_{10} 11101100_{2}=-19_{10}$
$-0_{10}: 00000000_{2}$ and $11111111_{2}$
- Two's complement - invert each bit and add 1
$-00010011_{2}=19_{10} 11101101_{2}=-19_{10}$
- Try to negate 0:
- $0_{10}=00000000_{2}$
- invert: $00000000_{2} \rightarrow 11111111_{2}$
- add 1: $11111111_{2}+00000001_{2}=00000000_{2}$


# Two's Complement Notation <br> Systems 

| Bit Pattern | Value Represented | Bit Pattern | Value <br> Represented |
| :---: | :---: | :---: | :---: |
| 011 | 3 | 0111 | 7 |
| 010 | 2 | 0110 | 6 |
| 001 | 1 | 0101 | 5 |
| 000 | 0 | 0100 | 4 |
| 111 | -1 | 0011 | 3 |
| 110 | -2 | 0010 | 2 |
| 101 | -3 | 0001 | 1 |
| 100 | -4 | 0000 | 0 |
|  |  | 1111 | -1 |
|  |  | 1110 | -2 |
|  |  | 1101 | -3 |
|  |  | 1100 | -4 |
|  |  | 1011 | -5 |
|  |  | 1010 | -6 |
|  |  | 1001 | -7 |
|  |  | 1000 | -8 |

## Addition



- Example: we are using 8 bit two's complement and have the problem 97-81.

| 97 | 0 | 1 | 1 | 0 | 0 | 0 | 0 | 1 |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| -81 | +1 | 0 | 1 | 0 | 1 | 1 | 1 | 1 |
| 16 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 |

## Summary: One's/Two's Complement

- Note that in sign magnitude and in one's complement there are two possible representations of 0 , and we can represent every integer $n$ with $-\left(2^{k}-1\right) \leq \mathrm{n} \leq 2^{k}-1$ with a k-bit field.
- With two's complement there is only one representation of 0 , and we can represent the integers $n$ with $-2^{k} \leq \mathrm{n} \leq 2^{k}-1$.
- The most significant advantage of two's complement is the fact that subtraction can be performed with addition circuitry.


## The Problem of Overflow

- Overflow: when a value to be represented falls outside the range of values that can be represented.
- An overflow is indicated if the addition of two positive values results in the pattern for a negative value or vice versa.
- Remember: small values can accumulate to produce large numbers.


## Fractions in Binary

- Each position is assigned a quantity of twice the size of the one to its right.
- $101.101=$ ?


[^0]
## Floating Point representation

- To represent a floating point number the number is divided into two parts: the integer and the fraction
- 3.1415 has integer 3 and fraction 0.1415
- Converting FP to binary:
- Convert integer part to binary
- Convert fraction part to binary
- Put a decimal point between the two numbers

Floating Point representation

- Assume 12 bits to represent the integer part and 4 bits to represent the fraction:
$71.3425=+1000111.0101$


## Normalization

## - To represent +1000111.0101

- Store sign, all bits, and the position of decimal point


## - Instead we use Normalization

-Move the decimal point so that there is only one 1 to the left of the decimal point.
-To indicate the original value of the number, multiply by $2^{\mathrm{e}}$ where e is the number of bits that the decimal point moved, positive for left movement, negative for right movement
-Store:
-The sign
-The exponent
-The mantissa


## Excess (or bias) Notation

- Alternative to two's complement used to store the exponent for floating point numbers.

| Bit Pattern Value <br> Represented <br> 111 3 <br> 110 2 <br> 101 1 <br> 100 0 <br> 011 -1 <br> 010 -2 <br> 001 -3 <br> 000 -4 |
| :--- |

With 3 bits we have excess 4 ( $=2^{3-1}$ )

## Excess (or bias) Notation

| Bit Pattern | Value <br> Represented |
| :---: | :---: |
| 1111 | 7 |
| 1110 | 6 |
| 1101 | 5 |
| 1100 | 4 |
| 1011 | 3 |
| 1010 | 2 |
| 1001 | 1 |
| 1000 | 0 |
| 0111 | -1 |
| 0110 | -2 |
| 0101 | -3 |
| 0100 | -4 |
| 0011 | -5 |
| 0010 | -6 |
| 0001 | -7 |
| 0000 | -8 |

With 4 bits we have excess 8 ( $=2^{4-1}$ )

With N bits we have excess $2^{\mathrm{N}-1}$

We add the magic number, $2^{\mathrm{N}-1}$, to the integer, change the result to binary, and add 0's (on the left) to make up N bits.

## FP Representation

- The standard IEEE representation for FP uses 32 bits for single-precision format:
- 1 bit sign
-8 bits exponent (excess $127=2^{8-1}-1$ )
- 23 bitsmantissa
- Representation. Store the:
- Sign as 0 (positive), 1 (negative)
- Exponent (power of 2) as excess127
- Mantissa as an unsigned integer


## Example: FP Representation

- Consider a 16 -bit representation with
- 1 bit sign
- 5 bits exponent (excess $16=2^{5-1}$ )
- 10 bitsmantissa
- $0 \quad 100111101100000$
- The sign 0 represents positive
- The exponent $10011=19_{10}$ represents 3
- The mantissa is 1101100000 ( 1 to left of . is not stored, it is understood)


## Coding the value 2 5/8




[^0]:    10.001
    $+100.000$
    111.001

