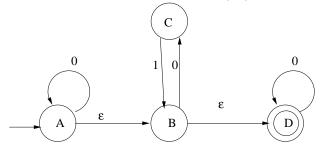
- 1. Give a Regular Expression and DFA for:  $L = \{x \in \{0, 1\}^* \mid x \text{ ends with } 1 \text{ and does not contain the substring } 00\}$
- 2. Give a RE for:  $L = \{0^i 1^j \mid i \text{ is even and } j \text{ is odd } \}$
- 3. Given the NFA for below for  $0^*(01)^*0^*$ , construct a DFA:

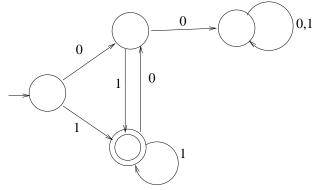


- 4. Give a RE and a DFA/NFA for the language of all strings over  $\{0, 1\}^*$  that do not end in 01.
- 5. Give a RE and a CFG for:  $L = \{x \in \{0, 1\}^* \mid x \text{ starts and ends with different symbols } \}$
- 6. Give a CFG for:  $L = \{x \in \{0, 1\}^* \mid \text{symbol at position } i \text{ is same as symbol at position } i+2 \text{ and } | x | \ge 2\}$
- 7. Give a CFG for the language of all non-palindromes over  $\{0, 1\}^*$ .
- 8. Give a CFG for:  $L = \{0^{i}1^{j}0^{k} \mid j > i+k\}$  So, 001111100 is in the string. Hint, the concatenation of two (or more) context-free languages is context-free.
- 9. Eliminate left recursion from:  $S \rightarrow Aa \mid b$  $A \rightarrow Ac \mid Sd \mid \varepsilon$
- 10. Give a CFG for  $L = \{a^i b^i c^i \mid i \ge 1\}$ .
- 11. Is this grammar ambiguous? If so, prove it and construct a non-ambiguous grammar that derives the same language.  $S \rightarrow aS \mid aSbS \mid c$

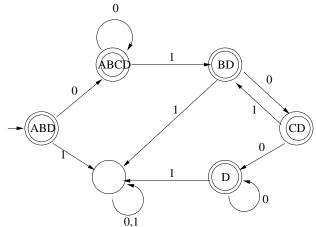
12. Assume that we have added a pointer type to decaf that can point to integers and booleans. We want to extend our type system (our attributed grammar) to handle these types. We have added a *pointer(t)* type to the type system to denote a pointer of type t. Complete the semantic action that propagates the type attribute for a pointer deference expression:  $E \to *E_1$  {: E.type = ???? :}

## Answers

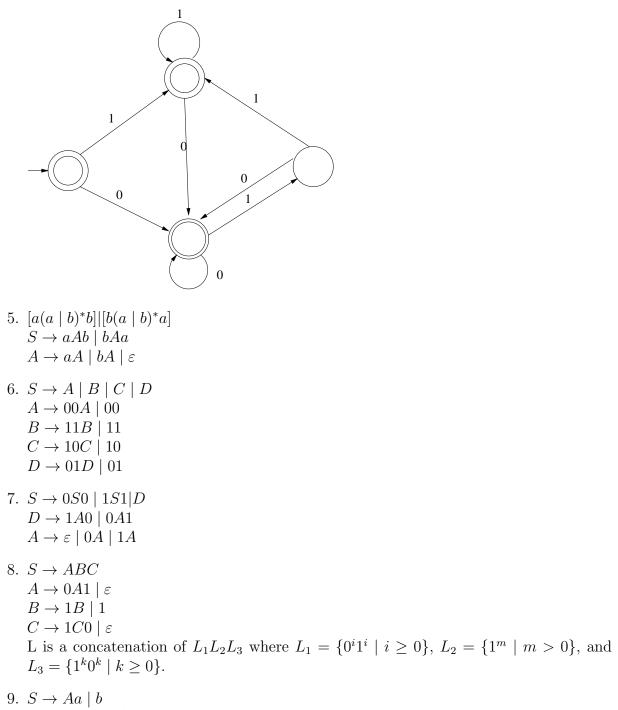
1.  $(1 \mid 01) +$ 



- 2.  $(00)^*1(11)^*$
- 3. DFA:



4.  $\varepsilon \mid 0 \mid 1 \mid (0 \mid 1)^* (11 \mid 00 \mid 10)$ 



$$\begin{array}{c} A \to bdA^1 \mid A^1 \\ A^1 \to cA^1 \mid adA^1 | \varepsilon \end{array}$$

- 10. Trick question, the language is not context-free. Sorry!
- 11. It is ambiguous! aacbc has two parse trees (not pictured, but you have to show the two parse trees to prove it is ambiguous).

Unambiguous grammar:  $S \rightarrow T \mid U$   $T \rightarrow aTbT \mid c$  $U \rightarrow aS \mid aTbU$ 

12.  $E \to *E_1$  {E.type := if  $E_1$ .type = pointer(t) then t else type\_error;}

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