# A Tale of 10 Bugs: Performance Engineering at VMware

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## Email thread from a colleague

"

. . .

"

Interestingly, as the number of <benchmark> threads decreased hostsPerThread var increases), the percentage of locktime spent in dbwrites also increases.

lots of threads (hostsPerThread = 4):

- ~28 % lock time spent under vdbWrite Connection
- ~16 % lock time spent under exec / commit.

Translation: Why is % lock time in DB increasing despite lighter load?



## Step 0: What the ?%##!\* is he talking about?



4. Client is notified of completion

Problem: With *lighter* load from client, %time spent in DB Locks *increases* 

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# **Step 1: Examine Lock Hold Time for Various Loads**



Latency per lock @ 128 hosts/thread < 4 hosts/thread (Expected...lighter load)

 $\rightarrow$ Original question: why is %DB increasing with lighter load?

 $\rightarrow$ Answer: DB latency dominates when overall latency is lower!

## **Step 2: Examine Contention Time for Various Loads**



Contention per lock @ 128 hosts/thread < 4 hosts/thread (OK…lighter load)</li>
→With lighter load, less overall contention time and higher % of time @ DB

- 1. Understand experimental setup (multi-tier setup)
- 2. Understand what is being measured (% time in DB lock)
- 3. Examine relevant data (lock latency)
- 4. Draw appropriate conclusion
  - Yes, % lock time in DB is higher with a lighter load
  - BUT, overall lock time is small with lighter load
  - Therefore, DB lock time (roughly constant) contributes more to lock latency

## **Case Studies in Performance Engineering @ VMware**

Lessons Learned



# Case Study #2: Garbage In, Garbage Out

Performance Chart Legend

### Customer wants to draw this chart:



Кеу	Object	Measurement	Rollup	Units	Latest	Maximum	Minimum	Average
	SpecJBB-esx1	CPU Usage in MHz	Average	MHz	4237	4472	78	421.883
	SpecJBB-esx1	CPU Ready	Summation	Millisecond	165	269	20	87.061
	0	CPU Ready	Summation	Millisecond	85	132	10	43.789
	1	CPU Ready	Summation	Millisecond	81	143	9	43.272
	0	CPU Usage in MHz	Average	MHz	2119	2233	40	210.611
	1	CPU Usage in MHz	Average	MHz	2102	2221	34	207.15

### **PowerCLI**

- CPU Usage for a VM for last hour:
- \$vm = Get-VM –Name "Foo"
- Get-Stat Entity \$vm Realtime Maxsample 180 Stat cpu.usagemhz.average
- Grab appropriate fields from output, use graphing program, etc.

# What Happens at Scale? Comparing PowerCLI and Java

Entities (cpu.usagemhz.average)	<b>PowerCLI</b> (Time in secs)	Java (Time in secs)
1 VM	9.2	14
6 VMs	11	14.5
39 VMs	101	16 Highly-tu
363 VMs	2580 (43 minutes)	50 Java St

A Naïve script that works for small environments may not be suitable for large environments

Translation: Garbage In, Garbage Out...but why?

# **PowerCLI vs. Java**

### **PowerCLI**

- Toolkit: meant for ease of use...hides details
- Similar to a shell script: facilitates quick prototyping
- Stateless

## Java

- Harder to use
- But...can use more advanced techniques (data structures, thread pools, etc.)



# What's going on behind the scenes?

- This is what is going on for each Get-Stat call in PowerCLI
- Retrieve PerformanceManager
- QueryPerfProviderSummary \$vm → Says what intervals are supported
- QueryAvailablePerfMetric \$vm → **Describes available metrics**
- QueryPerfCounter → Verbose description of counters
- Create PerfQuerySpec → Query specification to get the stats
- QueryPerf → Get stats

Bottom line: The PowerCLI toolkit spares you details...Easy to use!

# **Optimizing the Java Code**

### PowerCLI

Get VM ID

for each Get-Stat {

QueryAvailablePerfMetric();

QueryPerfCounter();

QueryPerfProviderSummary(); create PerfQuerySpec();

QueryPerf();



PowerCLI: 5 RPC calls per VM. Java: 1 RPC call per VM.

Further optimization not shown: Java allows more compact format

# Why Garbage In, Garbage out?

## **PowerCLI**

- Wrote a 'simple' but non-optimized script
- Did not utilize multi-threading (split up VM list, use muliple client queries)
- Did not realize output format is verbose
- Did not realize # of RPC calls is 5\*O(#VMs)

## Java

- Utilized multiple threads
- Understood what data was the same across VMs  $\rightarrow$  reduce redundant calls
- Utilized more compact output format (CSV vs. raw objects)
- Reduced # of RPC calls

(Think about assembly code vs. compiler-generated code)

# **Case Study #3: A Lesson in API Design**



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# The Problem: Remote Console Doesn't Show Up!

• Problem: could not start VM remote console in large environment

## Sequence of debugging

- Client folks: it's a server problem
- Server folks: it's a client problem
- Client folks: it's a 'vmrc' problem (vmrc = VMware Remote Console)
- VMRC folks: authentication? MKS tickets?
- Me: this is ridiculous...

### More Information: Start remote console for a single VM

- 50 Hosts, no problem
- 500 Hosts, no problem
- 1001 Hosts, PROBLEM!

## Debugging observations

- With < 1000 hosts...
  - Management server CPU and memory goes very high when client invoked
  - Console is dark until CPU and memory go down, then appears
- Look at server log file
  - Data retrieval call occurs before console appears (WHY???)
  - In failure cases, exception in serializer code
- Attach debugger
  - Exception is an out-of-memory exception
  - Exception is silently ignored (never returns to client)

# No Console: Isolating the Problem

## Problem

- VMRC creates a request to monitor host information (e.g., is CD-ROM attached)
- Request gets info on ALL hosts
- At 1001 hosts, we exceed 200MB buffer on server
- 200MB restriction only for old-style API clients

# Solution

- VMRC folks: do NOT create big request
- Server folks: fail correctly and emit better errors

### Lessons

- 1. Create APIs that are difficult to abuse, rather than easy to abuse
- 2. Teach clients how to use APIs
- 3. Make sure (internal) users have input about API design

### **Benchmark run**

- Build A: 100 ops/min.
- Build B: 50 ops/min.

### What was the difference?

- Build A: 32-bit executable on 64-bit hardware
- Build B: 64-bit executable on 64-bit hardware

Huh?

# 4 (b) xPerf

**Runs on Windows 2008** 

Sampling profiler (with other cool attributes)

**Records stack traces** 

**Give caller/callee information** 





CPU is mostly saturated (in 32-bit case, CPU is not saturated)

Stack	Weight -	/ % Weight	
	3,532,612.638	87.39	
E [Root]	3,532,455.583	87,38	
I- ntdll.dll!RtlUserThreadStart	3,508,983.437	86.80	
E  - ntdll.dll!ZwQueryVirtualMemory	13,762.103 969	0.34	
E  - ntkrnlmp.exe!KiSystemServiceCopyEnd	5,903.542 252	0.15	
E [- ntkrnlmp.exe!KiChainedDispatch	901.654 728	0.02	
E  - ntkrnlmp.exe!KiDpcInterrupt	321.063/214	0.01	
∃  - ntdll.dll! ?? ::FNODOBFM::`string'	258.007 162	0.01	
E  - ntdll.dll!RtlImageNtHeaderEx	226.071 622	0.01	
E  - MSVCR80.DLL!_RTtypeid	105.035 557	0.00	
The shall dispute the contract All strategies as	01 004 057	0.00	

Shows stacks originating from root

Shows 87% CPU used from 1 process

But this is just the thread start routine, where threads originate

Stack	Weight 👻 🕺 Weight
	3,532,612.638 87.39
🖂 [Root]	3,532,455.583 87.38
I- ntdll.dll!RtlUserThreadStart	3,508,983.437 86.80
kernel32.dll!BaseThreadInitThunk	3,508,983.437 86.80
🖂    - vpxd.exe!Win32ThreadMain	2,270,619.910 56.17
vpxd.exe!VpxdThread::ThreadFunc	2,270,619.910 56.17
I   - vpxd.exe!VpxLroList::ThreadMainEntry	2,045,997.133 50.61
🖂       - vpxd.exe!VpxLRO::LroMain	2,041,093.903 50.49
vpxd.exe!VpxActivationLRO::Inv	okeA 1,571,975.867 38.89
🖂          - vmomi.dll!Vmomi::ManagedM	ethod 1,518,822.158 37.57
🗄           - vmomi.dll!sVmodlQueryPr	opert 770,053.167 788 19.05
🖃           - types.dll!sVimVirtualMach	ineDi 727,687.116 523 / 18.00
□               - vpxd.exe!VpxdMoVm	::Po 672,630.405 844 / 16.64

From Root, most of the samples are from this call stack Most popular stack, but is this the problem?

# 4(f) Perils of Sampling Profilers, Part 2

### Most-common trace: not necessarily where time is spent



# 4(g) The Caller View

## Look at Callers for various routines in stacks

Callers	Weight 👻	% Weight	
ntdll.dll!ZwQueryVirtualMemory	3,123,003.752	77.26	
I - ntdll.dll! ?? ::FNODOBFM::`string'	3,109,241.648	76.92	
∃   - MSVCR80.DLL!_RTDynamicCast	1,579,519.929	39.07	
∃   - MSVCR80.DLL!_RTtypeid	1,529,451.706	137.84	
	257,005 987	0.01	
	13,006,609	0.00	
	13,762.103 969	0.34	
Intell.dll!ZwQueryVirtualMemory	2.004 444	0.00	

Not called a lot from root, however...

Called from few places and takes 77% CPU!

RTtypeid?

Callers		Weight 🔻	% Weight	
E	MSVCR80.DLL!_RTtypeid	1,583,683.660	39.18	
Ŧ	- vmacore.dll!Vmacore::ObjectImpl::IncRef	828,866.361 590	20.50	
Ŧ	- vmacore.dll!Vmacore::ObjectImpl::DecRef	725,473.308 729	17.95	
+	- MSVCR80.DLL!_RTDynamicCast	28,925.004 534	0.72	
+	<pre> - vpxd.exe!ManagedObjectMapper::operator()</pre>	283.942 793	0.01	
+	- [Root]	105.085 557	0.00	
+	- vpxd.exe!VpxLRO::GetStatsContext	18.005 567	0.00	
Ŧ	- vpxd.exe!DrmModule::SnapshotDomain	10.999 304	0.00	
+	- vmacore.dll!Vmacore::PrintFormatter::FormatException	1.002 432	0.00	
Ŧ	- vmacore.dll!Vmacore::PrintFormatter::FormatException	1.002 4	132	

Hmm. RTtypeid is used in figuring out C++ type 39% of overall CPU?

IncRef and DecRef are main callers

# 4(i) The Offending Code

```
void
ObjectImpl::IncRef()
{
    if (_refCount.ReadInc() == 0) {
        const type_info& tinfo = typeid(*this);
        FirstIncRef(tinfo);
    }
    ...
}
Dynamic cast...needs run-time type info (RTTI)
```

**RTTI** has pointers in it

# 4(j) But Why is 64-bit slower than 32-bit?

## Runtime type info (RTTI) has a bunch of pointers

- 32-bit: pointers are raw 32-bit pointers
- 64-bit
  - Pointers are 32-bit offsets
  - Offsets must be added to base addr of DLL/EXE in which RTTI resides
  - Result is a true 64-bit pointer

## But wait...why is addition slow?

# 4(k) Why Is Addition Slow? Well, it isn't...

## Addition isn't slow, but...

## Determining module base address can be slow

- To find base address, RTtypeid calls RtIPcToFileHeader
- RtIPcToFileHeader grabs loader lock, walks list of loaded modules to find RTTI data
- This can be slow
- N.B.: This is why we see calls to zwQueryVirtualMemory

For more info: http://blogs.msdn.com/junfeng/archive/2006/10/17/dynamic-cast-isslow-in-x64.aspx

# 4(I) What Did We Learn?

RtTypeId is called from a bunch of places

RtTypeId is not, however, called from Root too often

RtTypeId is small and fast: not main contributor in most stacks (except IncRef and DecRef)

Lots of little calls add up

Caller view was important here!

(btw: 2 solutions:

- 1. Statically compute base addr and cache
- 2. Use latest runtime library, which avoids RtIToPcFileHeader)

Lesson: Little things (32-bit vs. 64-bit) may matter...don't discriminate!

# **Case Study #5: Memory Usage Woes**

### Why is excessive memory usage a problem?

- Can slow down application if paging is induced
- May cause application to crash (if you exceed per-process limit...2GB in 32-bit Windows)

### Memory leak vs. memory accumulation

- Leak: memory was allocated, not live anymore (dangling reference)
- Accumulation: pointer exists to data, but data not used anymore (a logical leak)

# **Tools for Analyzing Memory Usage**

## Windows:

• Purify, GlowCode, Memory Validator, malloc hooks and heap dump utilities from Microsoft, etc.

# Linux:

 Valgrind, malloc hooks from Google (example: http://googperftools.sourceforge.net/), etc.

## Basic idea:

- Hook calls to malloc
- Figure out liveness of pointers (do you leave scope without free()?)
- But...can be unusably slow if you do a lot of allocations!



## void bar() {

foo(); }

# void foo() {

```
char *p = malloc(24);
```

<do some computation>

return; /\* memory pointed at by p is never freed \*/

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}

# **Memory Analysis**

# Easing memory allocation in C++: use reference-counted objects instead of "naked" pointers

- Each use of an item increments a reference count
- When no references exist, delete the item
- Does not solve memory accumulation problem

## Server application runs out of memory after several hours

## Use Purify (on a much smaller setup):

- Leak not detected because data was assigned to a reference
- Instead, examine memory in use
  - Do 100 iterations of an operation
  - See 6400B of allocations for an item (100 64B allocations)
  - Code inspection revealed that item was actually not used anymore...a "logical" leak
  - (i.e., there was a free(), but it was never called because the item was thought to be in use)

Lesson:

If an effect is small, find ways to magnify it.



# **Case Study #6: Another Memory Analysis Problem**

User complains that server is getting slower and slower

**CPU/network/disk not saturated** 

Memory, however, is increasing dramatically

**Eventually, system crashes** 



# Looking at Memory Usage: Perfmon in Windows

## Chart of "Private Bytes" for a process vs. time



Server is functioning fine, but memory is growing really fast. This could lead to a crash. Let's investigate...

## **Profiling Reference-counted Objects**



# **Customized Profiling: Pros and Cons**

## Advantages of our customized profiler:

- Tailored to our application
- Can be made very fast
- Can be run in production environments

## **Disadvantages:**

- Requires code recompilation (then again, so does Purify)
- Specific to this application (code must be refactored for use in other apps)
- Only counts ref-counted objects: what about C code? What about non-refcounted objects?

Lesson: Memory profiling is critical.

Sad Reality: Sometimes, commercial tools don't work at scale

 $\rightarrow$ You may have to write your own

# Case Study #7: How well do you understand networking?

### User issues a request to perform an operation on a VM

- Setup A: Client/Server version 1 to host version 1: 8s
- Setup B: Client/Server version 2 to host version 1: 16s
- Consistent, repeatable difference
- Regression when using new code to talk to older host!

## Step 1: Log everywhere

- Client-imposed latency: same in both cases
- Server-imposd latency: same
- Host imposed-latency: extra 8s in Setup B  $\rightarrow$  Focus on the host



# **Networking Issue: Analyzing the host**

## Step 2: More logging (standard tools aren't available on host)

- Narrow down the time...
  - Agent <-> HAL, Setup A: 10ms per call
  - Agent <-> HAL, Setup B: 200ms per call
  - Wow!

## **Step 3: Examine configuration**

- Setup A: named pipe between Agent and HAL
- Setup B: TCP/IP connection between Agent and HAL

# **Networking Issue: Resolution**

# Step 4: Solution (intuition by developer)

- Named pipe communication, setup A: 10ms
- TCP/IP communication, setup B: 200ms
- Why? Nagle algorithm on socket connection
  - On a TCP socket, wait for more data before sending packets
  - Can be disabled through TCP\_NODELAY option

## Step 5: Result

- Use TCP\_NODELAY, both have same performance
- Eventually use a cache to avoid interprocess communication

## Lesson?

- "Little" changes can mean a lot
- Client/server code: understand the client/server interaction!

# **Case Study #8: Correctness Impacts Performance**

### Trying to Power on a VM

- Sometimes, powering on VM would take 5 seconds
- Other times, powering on VM would take 5 minutes!

### Where to begin?

• Powering on a VM requires disk activity on host → Check disk metrics for host

# Examining Disk Latencies...



→Chart shows highest disk latency for each 5-minute period
→Max Disk Latencies range from 100ms to 1100ms...very high! Why?

# **High Disk Latency: Mystery Solved**



Host events: disk has connectivity issues  $\rightarrow$  high latencies!

Bottom line: correctness issue (bad disk controller) impacts performance!

# **Prelude to Case Studies 9 & 10: CPU Scheduling for VMs**



Run (accumulating used time) Ready (wants to run, no physical CPU available) Wait: blocked on I/O or voluntarily descheduled

# Case Study 9: "But it's only a small probe VM..."



vSphere communicates with DB

Probe VM monitors vSphere-to-DB traffic

The more traffic, the more work done by Probe VM

User Complaint: vSphere VM is suddenly very unresponsive

# **CPU Usage vs. Time for DB and Probe VM**



DB VM ready time goes from 12.5% when idle to ~20% when user busy

DB ready time increases because Probe VM is busy

Probe VM takes CPU away from DB VM  $\rightarrow$  user responsiveness suffers

# **Case Study #10: What Does This Metric Mean?**

Problem

- Customer Performs a Load Test: keeps attaching clients to a server
- At some point, CPU is NOT saturated, but latency starts to degrade
- At some point, client is unusable
- Why?

## "Oh yeah, it's a disk problem..."



## Uh-oh! Disk Latencies go over a cliff!

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**Problem:** 

Yes, Disk Latency gets worse at 4pm. (btw...due to swapping)

However, Application latency gets worse at 3:30pm!

What's going on from 3:30pm to 4pm?



# Looking at a different chart...

ID	GID	NAME	NWLD	%USED	∜RUN	\$SYS	%WAIT	%RDY	%IDLE	*
1	1	idle	16	111.77	563.57	0.00	0.00	800.00	0.00	
2	2	system	7	0.01	0.02	0.00	700.00	0.00	0.00	
6	6	helper	73	0.15	0.25	0.00	7300.00	0.35	0.00	
7	7	drivers	9	0.00	0.01	0.00	900.00	0.01	0.00	
8	8	vmotion	4	0.00	0.00	0.00	400.00	0.00	0.00	
10	10	console	2	6.45	10.04	0.01	186.59	4.53	86.00	
15	15	vmkapimod	5	0.19	0.28	0.00	500.00	0.00	0.00	
17	17	FΤ	1	0.00	0.00	0.00	100.00	0.00	0.00	
18	18	vobd.4279	8	0.00	0.00	0.00	800.00	0.00	0.00	
19	19	net-cdp.4287	1	0.00	0.00	0.00	100.00	0.00	0.00	
20	20	vmware-vmkauthd	1	U.00	0.29	0.00	100.00	0.00	0.00	
68	68	vm1	5	183.21	248.36	0.28	232.84	18.64	25.13	
69	69	vm2	5	152.17	212.16	0.37	284.93	5.26	77.26	
70	70	vm3	5	126.52	194.13	0.17	302.64	5.64	93.50	
71	71	vm4	5	146.25	219.30	0.21	270.28	11.54	64.30	

%Used? %Run? What's the difference?

%used: normalized to base clock frequency

Lesson: understand

your metrics!

%run: normalized to clock frequency while VM is running...

%run > %used: Power Management is kicking in...

In this case, turn off power management -> latency problems go away

# The 10 Performance Issues I Mentioned

## 1. DB Lock % increase with decreasing load

• Be careful when you draw conclusions...

## 2. PowerCLI vs. Java

• Garbage-In, Garbage-Out: scalable solutions require careful design

## 3. Remote Console Issues

• Create APIs that are easy to use and difficult to abuse

## 4. 32-bit vs. 64-bit

• A small change can make a HUGE difference

# 5. "Logical" leak

- Just because you do "new/delete," doesn't mean memory won't grow (btw., Java doesn't save you!)
- Exaggerate a problem to make it easier to find the root cause



# The 10 Performance Issues I Mentioned

## 6. Slow memory growth until crash

• Sometimes you need customized profilers

# 7. Nagling

• Understand client/server interactions

# 8. Disk Latency

Correctness Impacts Performance

# 9. Probe VM activity hurting performance of other VMs

• Understand the Impact of Monitoring

# **10. Power Management affecting Performance**

• Understand your metrics & consider the whole system

Avoid assumptions! (see #10)

# Understand the ENTIRE SYSTEM

- Your code
- Other people's code
- Hardware

## Be persistent and thorough

- Look at tons of metrics
- Look at behavior when things work as well as when they don't work



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