

6.252 NONLINEAR PROGRAMMING

LECTURE 19: DUALITY THEOREMS

LECTURE OUTLINE

- Duality and L-multipliers (continued)
- Consider the problem

minimize $f(x)$

subject to $x \in X, \quad g_j(x) \leq 0, \quad j = 1, \dots, r,$

assuming $-\infty < f^* < \infty$.

- μ^* is a Lagrange multiplier if $\mu^* \geq 0$ and $f^* = \inf_{x \in X} L(x, \mu^*)$.

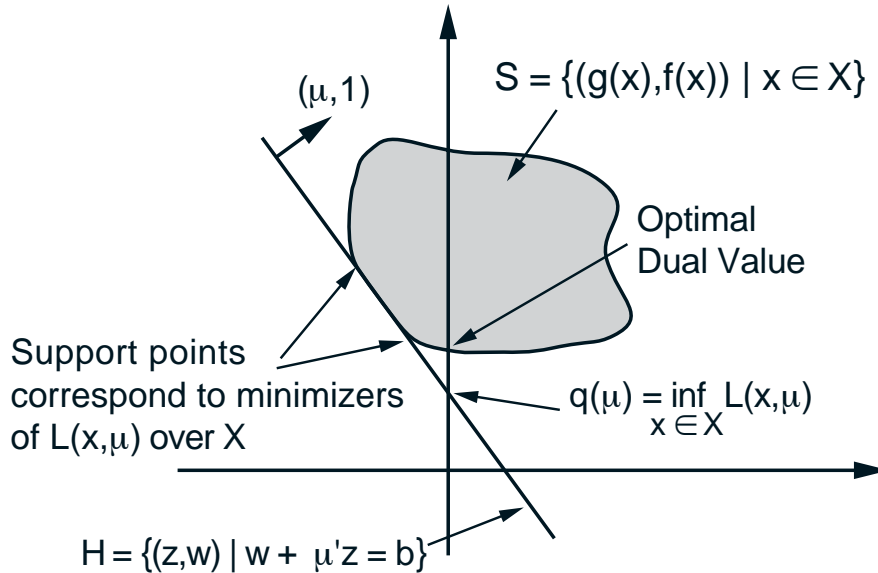
- The *dual problem* is

maximize $q(\mu)$

subject to $\mu \geq 0,$

where q is the dual function $q(\mu) = \inf_{x \in X} L(x, \mu)$.

DUAL OPTIMALITY



- Weak Duality Theorem: $q^* \leq f^*$.
- Lagrange Multipliers and Dual Optimal Solutions:
 - (a) If there is no duality gap, the set of Lagrange multipliers is equal to the set of optimal dual solutions.
 - (b) If there is a duality gap, the set of Lagrange multipliers is empty.

DUALITY PROPERTIES

- **Optimality Conditions:** (x^*, μ^*) is an optimal solution-Lagrange multiplier pair if and only if

$$x^* \in X, \quad g(x^*) \leq 0, \quad (\text{Primal Feasibility}),$$

$$\mu^* \geq 0, \quad (\text{Dual Feasibility}),$$

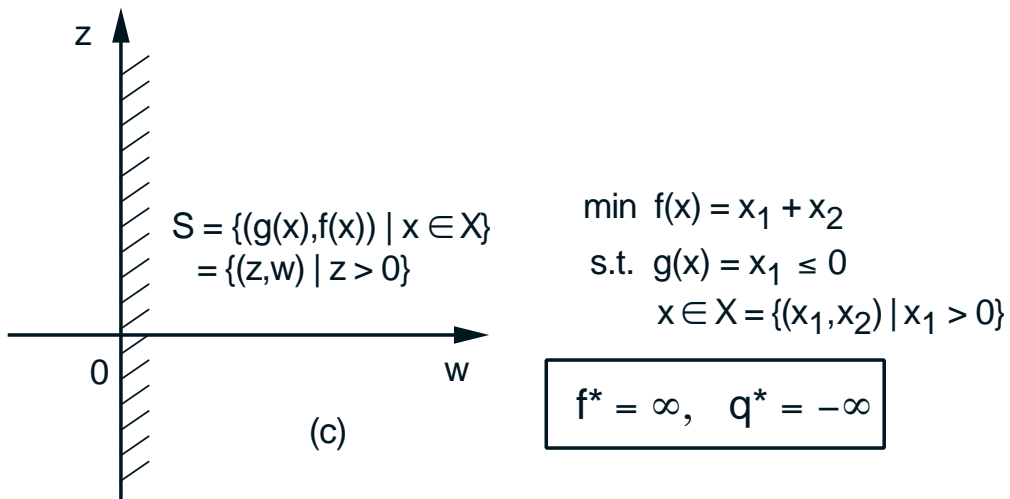
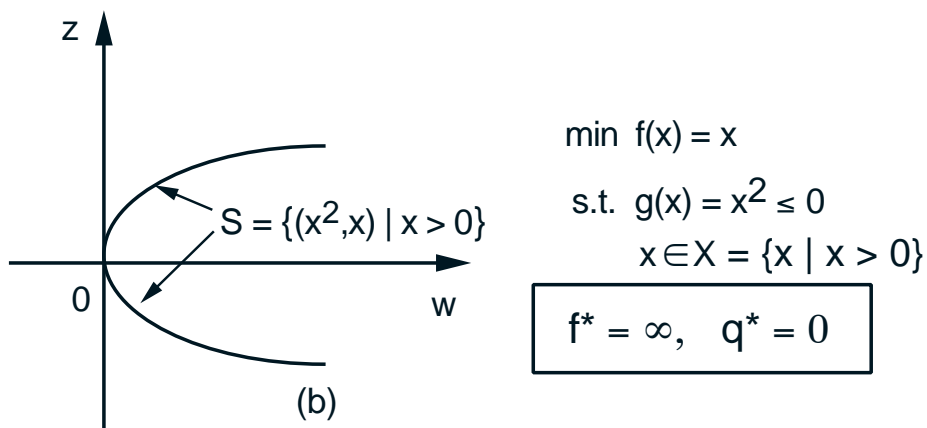
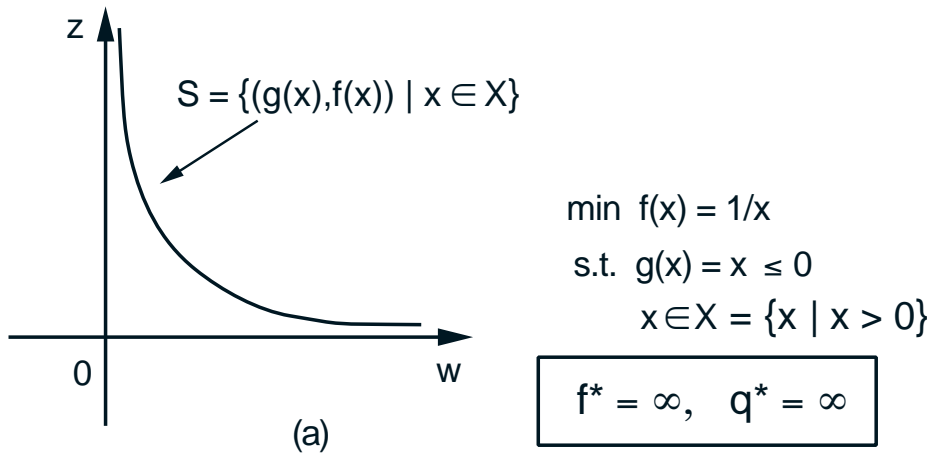
$$x^* = \arg \min_{x \in X} L(x, \mu^*), \quad (\text{Lagrangian Optimality}),$$

$$\mu_j^* g_j(x^*) = 0, \quad j = 1, \dots, r, \quad (\text{Compl. Slackness}).$$

- **Saddle Point Theorem:** (x^*, μ^*) is an optimal solution-Lagrange multiplier pair if and only if $x^* \in X$, $\mu^* \geq 0$, and (x^*, μ^*) is a saddle point of the Lagrangian, in the sense that

$$L(x^*, \mu) \leq L(x^*, \mu^*) \leq L(x, \mu^*), \quad \forall x \in X, \mu \geq 0.$$

INFEASIBLE AND UNBOUNDED PROBLEMS



EXTENSIONS AND APPLICATIONS

- Equality constraints $h_i(x) = 0$, $i = 1, \dots, m$, can be converted into the two inequality constraints

$$h_i(x) \leq 0, \quad -h_i(x) \leq 0.$$

- Separable problems:

$$\begin{aligned} &\text{minimize} && \sum_{i=1}^m f_i(x_i) \\ &\text{subject to} && \sum_{i=1}^m g_{ij}(x_i) \leq 0, \quad j = 1, \dots, r, \\ &&& x_i \in X_i, \quad i = 1, \dots, m. \end{aligned}$$

- Separable problem with a single constraint:

$$\begin{aligned} &\text{minimize} && \sum_{i=1}^n f_i(x_i) \\ &\text{subject to} && \sum_{i=1}^n x_i \geq A, \quad \alpha_i \leq x_i \leq \beta_i, \quad \forall i. \end{aligned}$$

DUALITY THEOREM I FOR CONVEX PROBLEMS

- Strong Duality Theorem - Linear Constraints:
Assume that the problem

minimize $f(x)$

subject to $x \in X$, $a'_i x - b_i = 0$, $i = 1, \dots, m$,

$e'_j x - d_j \leq 0$, $j = 1, \dots, r$,

is feasible and its optimal value f^* is finite. Let also f be convex over \mathbb{R}^n and let X be polyhedral. Then there exists at least one Lagrange multiplier and there is no duality gap.

- Proof Issues
- Application to Linear Programming

COUNTEREXAMPLE

- A Convex Problem with a Duality Gap: Consider the two-dimensional problem

minimize $f(x)$

subject to $x_1 = 0$, $x \in X = \{x \mid x \geq 0\}$,

where

$$f(x) = e^{-\sqrt{x_1 x_2}}, \quad \forall x \in X,$$

and $f(x)$ is arbitrarily defined for $x \notin X$.

- f is convex over X (its Hessian is positive definite in the interior of X), and $f^* = 1$.
- Also, for all $\mu \geq 0$ we have

$$q(\mu) = \inf_{x \geq 0} \left\{ e^{-\sqrt{x_1 x_2}} + \mu x_1 \right\} = 0,$$

since the expression in braces is nonnegative for $x \geq 0$ and can approach zero by taking $x_1 \rightarrow 0$ and $x_1 x_2 \rightarrow \infty$. It follows that $q^* = 0$.

DUALITY THEOREM II FOR CONVEX PROBLEMS

- Consider the problem

minimize $f(x)$

subject to $x \in X, \quad g_j(x) \leq 0, \quad j = 1, \dots, r.$

- Assume that X is convex and the functions $f : \mathcal{R}^n \mapsto \mathcal{R}, g_j : \mathcal{R}^n \mapsto \mathcal{R}$ are convex over X . Furthermore, the optimal value f^* is finite and there exists a vector $\bar{x} \in X$ such that

$$g_j(\bar{x}) < 0, \quad \forall j = 1, \dots, r.$$

- Strong Duality Theorem: There exists at least one Lagrange multiplier and there is no duality gap.
- Extension to linear equality constraints.