# LECTURE 20

# LECTURE OUTLINE

- Proximal methods
- Review of Proximal Minimization
- Proximal cutting plane algorithm
- Bundle methods
- Augmented Lagrangian Methods
- Dual Proximal Minimization Algorithm

• Method relationships to be established:



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# **RECALL PROXIMAL MINIMIZATION**



• Minimizes closed convex proper f:

$$x_{k+1} = \arg\min_{x \in \Re^n} \left\{ f(x) + \frac{1}{2c_k} \|x - x_k\|^2 \right\}$$

where  $x_0$  is an arbitrary starting point, and  $\{c_k\}$  is a positive parameter sequence.

• We have  $f(x_k) \to f^*$ . Also  $x_k \to$  some minimizer of f, provided one exists.

• Finite convergence for polyhedral f.

• Each iteration can be viewed in terms of Fenchel duality.

## **PROXIMAL/BUNDLE METHODS**

• Replace f with a cutting plane approx. and/or change quadratic regularization more conservatively.

• A general form:

$$x_{k+1} \in \arg\min_{x \in X} \{F_k(x) + p_k(x)\}$$

 $F_k(x) = \max\left\{f(x_0) + (x - x_0)'g_0, \dots, f(x_k) + (x - x_k)'g_k\right\}$  $p_k(x) = \frac{1}{2c_k} \|x - y_k\|^2$ 

where  $c_k$  is a positive scalar parameter.

• We refer to  $p_k(x)$  as the proximal term, and to its center  $y_k$  as the proximal center.



Change  $y_k$  in different ways => different methods.

# **PROXIMAL CUTTING PLANE METHODS**

• Keeps moving the proximal center at each iteration  $(y_k = x_k)$ 

• Same as proximal minimization algorithm, but f is replaced by a cutting plane approximation  $F_k$ :

$$x_{k+1} \in \arg\min_{x \in X} \left\{ F_k(x) + \frac{1}{2c_k} \|x - x_k\|^2 \right\}$$

where

$$F_k(x) = \max\{f(x_0) + (x - x_0)'g_0, \dots, f(x_k) + (x - x_k)'g_k\}$$

• Drawbacks:

- (a) **Stability issue:** For large enough  $c_k$  and polyhedral X,  $x_{k+1}$  is the exact minimum of  $F_k$  over X in a single minimization, so it is identical to the ordinary cutting plane method. For small  $c_k$  convergence is slow.
- (b) The number of subgradients used in  $F_k$ may become very large; the quadratic program may become very time-consuming.

• These drawbacks motivate algorithmic variants, called *bundle methods*.

#### **BUNDLE METHODS**

• Allow a proximal center  $y_k \neq x_k$ :  $x_{k+1} \in \arg\min_{x \in X} \left\{ F_k(x) + p_k(x) \right\}$   $F_k(x) = \max \left\{ f(x_0) + (x - x_0)'g_0, \dots, f(x_k) + (x - x_k)'g_k \right\}$ 

$$p_k(x) = \frac{1}{2c_k} ||x - y_k||^2$$

• Null/Serious test for changing  $y_k$ : For some fixed  $\beta \in (0, 1)$ 

$$y_{k+1} = \begin{cases} x_{k+1} & \text{if } f(y_k) - f(x_{k+1}) \ge \beta \delta_k, \\ y_k & \text{if } f(y_k) - f(x_{k+1}) < \beta \delta_k, \end{cases}$$
$$\delta_k = f(y_k) - \left( F_k(x_{k+1}) + p_k(x_{k+1}) \right) > 0$$



### **REVIEW OF FENCHEL DUALITY**

• Consider the problem

minimize  $f_1(x) + f_2(x)$ subject to  $x \in \Re^n$ ,

where  $f_1$  and  $f_2$  are closed proper convex.

#### • Duality Theorem:

- (a) If  $f^*$  is finite and  $\operatorname{ri}(\operatorname{dom}(f_1)) \cap \operatorname{ri}(\operatorname{dom}(f_2)) \neq \emptyset$ , then strong duality holds and there exists at least one dual optimal solution.
- (b) Strong duality holds, and  $(x^*, \lambda^*)$  is a primal and dual optimal solution pair if and only if

$$x^* \in \arg\min_{x \in \Re^n} \left\{ f_1(x) - x'\lambda^* \right\}, \ x^* \in \arg\min_{x \in \Re^n} \left\{ f_2(x) + x'\lambda^* \right\}$$

• By Fenchel inequality, the last condition is equivalent to

$$\lambda^* \in \partial f_1(x^*)$$
 [or equivalently  $x^* \in \partial f_1^*(\lambda^*)$ ]

and

 $-\lambda^* \in \partial f_2(x^*)$  [or equivalently  $x^* \in \partial f_2^*(-\lambda^*)$ ]

## **GEOMETRIC INTERPRETATION**



• When  $f_1$  and/or  $f_2$  are differentiable, the optimality condition is equivalent to

$$\lambda^* = \nabla f_1(x^*)$$
 and/or  $\lambda^* = -\nabla f_2(x^*)$ 

### DUAL PROXIMAL MINIMIZATION

• The proximal iteration can be written in the Fenchel form:  $\min_x \{f_1(x) + f_2(x)\}$  with

$$f_1(x) = f(x), \qquad f_2(x) = \frac{1}{2c_k} ||x - x_k||^2$$

• The Fenchel dual is

minimize  $f_1^{\star}(\lambda) + f_2^{\star}(-\lambda)$ subject to  $\lambda \in \Re^n$ 

• We have  $f_2^{\star}(-\lambda) = -x'_k \lambda + \frac{c_k}{2} \|\lambda\|^2$ , so the dual problem is

minimize  $f^{\star}(\lambda) - x'_k \lambda + \frac{c_k}{2} \|\lambda\|^2$ subject to  $\lambda \in \Re^n$ 

where  $f^*$  is the conjugate of f.

•  $f_2$  is real-valued, so no duality gap.

• Both primal and dual problems have a unique solution, since they involve a closed, strictly convex, and coercive cost function.

#### DUAL PROXIMAL ALGORITHM

• Can solve the Fenchel-dual problem instead of the primal at each iteration:

$$\lambda_{k+1} = \arg\min_{\lambda \in \Re^n} \left\{ f^*(\lambda) - x'_k \lambda + \frac{c_k}{2} \|\lambda\|^2 \right\}$$
(1)

• Lagragian optimality conditions:

$$x_{k+1} \in \arg \max_{x \in \Re^n} \left\{ x' \lambda_{k+1} - f(x) \right\}$$

$$x_{k+1} = \arg\min_{x \in \Re^n} \left\{ x' \lambda_{k+1} + \frac{1}{2c_k} \|x - x_k\|^2 \right\}$$

or equivalently,

$$\lambda_{k+1} \in \partial f(x_{k+1}), \qquad \lambda_{k+1} = \frac{x_k - x_{k+1}}{c_k}$$

• **Dual algorithm:** At iteration k, obtain  $\lambda_{k+1}$  from the dual proximal minimization (1) and set

$$x_{k+1} = x_k - c_k \lambda_{k+1}$$

• As  $x_k$  converges to a primal optimal solution  $x^*$ , the dual sequence  $\lambda_k$  converges to 0 (a subgradient of f at  $x^*$ ).

# VISUALIZATION



### • The primal and dual implementations are mathematically equivalent and generate identical sequences $\{x_k\}$ .

• Which one is preferable depends on whether f or its conjugate  $f^*$  has more convenient structure.

• Special case: When -f is the dual function of the constrained minimization  $\min_{g(x) \leq 0} F(x)$ , the dual algorithm is equivalent to an important general purpose algorithm: the Augmented Lagrangian method.

• This method aims to find a subgradient of the primal function  $p(u) = \min_{g(x) \le u} F(x)$  at u = 0 (i.e., a dual optimal solution).

### AUGMENTED LAGRANGIAN METHOD

• Consider the convex constrained problem

 $\begin{array}{ll} \text{minimize} & f(x) \\ \text{subject to} & x \in X, \quad Ex = d \end{array}$ 

• Primal and dual functions:

$$p(u) = \inf_{\substack{x \in X \\ Ex - d = u}} f(x), \quad q(\mu) = \inf_{x \in X} \left\{ f(x) + \mu'(Ex - d) \right\}$$

- Assume p: closed, so (q, p) are "conjugate" pair.
- Proximal algorithms for maximizing q:

$$\mu_{k+1} = \arg \max_{\mu \in \Re^m} \left\{ q(\mu) - \frac{1}{2c_k} \|\mu - \mu_k\|^2 \right\}$$
$$u_{k+1} = \arg \min_{u \in \Re^m} \left\{ p(u) + \mu'_k u + \frac{c_k}{2} \|u\|^2 \right\}$$

Dual update:  $\mu_{k+1} = \mu_k + c_k u_{k+1}$ 

• Implementation:

 $u_{k+1} = Ex_{k+1} - d,$   $x_{k+1} \in \arg\min_{x \in X} L_{c_k}(x, \mu_k)$ 

where  $L_c$  is the Augmented Lagrangian function

$$L_c(x,\mu) = f(x) + \mu'(Ex - d) + \frac{c}{2} ||Ex - d||^2$$

### **GRADIENT INTERPRETATION**

• Back to the dual proximal algorithm and the dual update  $\lambda_{k+1} = \frac{x_k - x_{k+1}}{c_k}$ 

• **Proposition:**  $\lambda_{k+1}$  can be viewed as a gradient:

$$\lambda_{k+1} = \frac{x_k - x_{k+1}}{c_k} = \nabla \phi_{c_k}(x_k),$$

where

$$\phi_c(z) = \inf_{x \in \Re^n} \left\{ f(x) + \frac{1}{2c} \|x - z\|^2 \right\}$$



• So the dual update  $x_{k+1} = x_k - c_k \lambda_{k+1}$  can be viewed as a gradient iteration for minimizing  $\phi_c(z)$  (which has the same minima as f).

• The gradient is calculated by the dual proximal minimization. Possibilities for faster methods (e.g., Newton, Quasi-Newton). Useful in augmented Lagrangian methods.

### **PROXIMAL LINEAR APPROXIMATION**

• Convex problem: Min  $f : \Re^n \mapsto \Re$  over X.

• Proximal outer linearization method: Same as proximal minimization algorithm, but f is replaced by a cutting plane approximation  $F_k$ :

$$x_{k+1} \in \arg\min_{x \in \Re^n} \left\{ F_k(x) + \frac{1}{2c_k} \|x - x_k\|^2 \right\}$$

$$\lambda_{k+1} = \frac{x_k - x_{k+1}}{c_k}$$

where  $g_i \in \partial f(x_i)$  for  $i \leq k$  and

$$F_k(x) = \max\{f(x_0) + (x - x_0)'g_0, \dots, f(x_k) + (x - x_k)'g_k\} + \delta_X(x)$$

• Proximal Inner Linearization Method (Dual proximal implementation): Let  $F_k^*$  be the conjugate of  $F_k$ . Set

$$\lambda_{k+1} \in \arg\min_{\lambda \in \Re^n} \left\{ F_k^{\star}(\lambda) - x_k^{\prime}\lambda + \frac{c_k}{2} \|\lambda\|^2 \right\}$$
$$x_{k+1} = x_k - c_k \lambda_{k+1}$$

Obtain  $g_{k+1} \in \partial f(x_{k+1})$ , either directly or via

$$g_{k+1} \in \arg \max_{\lambda \in \Re^n} \left\{ x'_{k+1}\lambda - f^*(\lambda) \right\}$$

• Add  $g_{k+1}$  to the outer linearization, or  $x_{k+1}$  to the inner linearization, and continue.

# **PROXIMAL INNER LINEARIZATION**

• It is a mathematical equivalent dual to the outer linearization method.



• Here we use the conjugacy relation between outer and inner linearization.

• Versions of these methods where the proximal center is changed only after some "algorithmic progress" is made:

- The outer linearization version is the (standard) bundle method.
- The inner linearization version is an **inner approximation version of a bundle method**.

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