

6.896  
2/4/04  
L1.1

## 6.896 Theory of Parallel Hardware.

Administrivia «Mention webpage bug»  
Overview «Discuss models»

### Systolic computation

E.g. linear array I/O — ○ — ○ — ○ — ○ — ... — ○

"Fixed-connection" network.

1. Underlying graph fixed.
2. Local communication only
3. I/O location restricted.

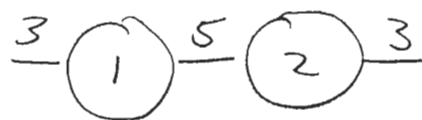
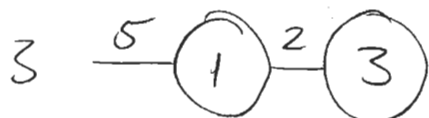
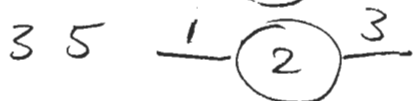
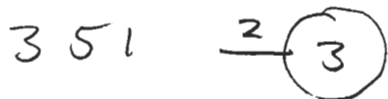
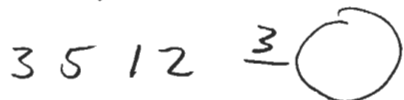
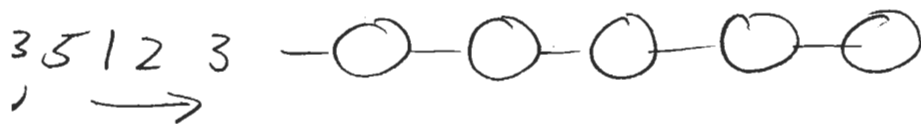
At each step of a globally synchronous clock, each processor.

1. receives inputs from neighbors (or I/O)
2. inspects local memory
3. performs local computation
4. updates local memory
5. generates outputs for neighbors.

### Example: Sorting

- Accept left input
- Compare input to stored value
- Store smaller value
- Output bigger # to right.

6.896  
2/4/04  
L1.2



⋮

Correctness: induction

$N$  inputs. How many steps?  $2N = \Theta(N)$

«Discuss outputting of values.»

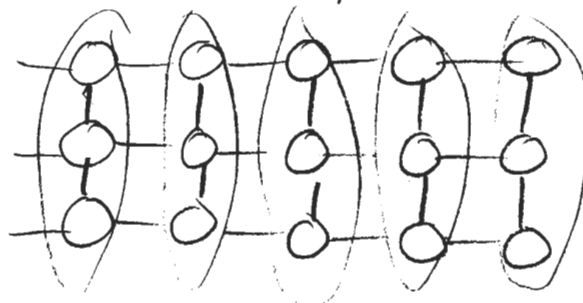
Total time =  $3N$

Sorting in the bit model (vs. word model)

• One processor per bit.  
 $N$   $k$ -bit #'s

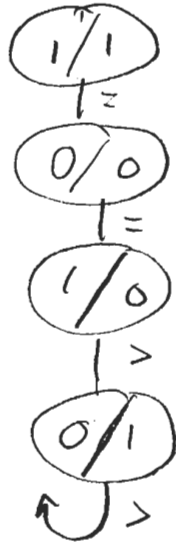
3	5	1	2	3
0	1	0	0	0
1	0	0	1	1
1	1	1	0	1

$N \times k$  array



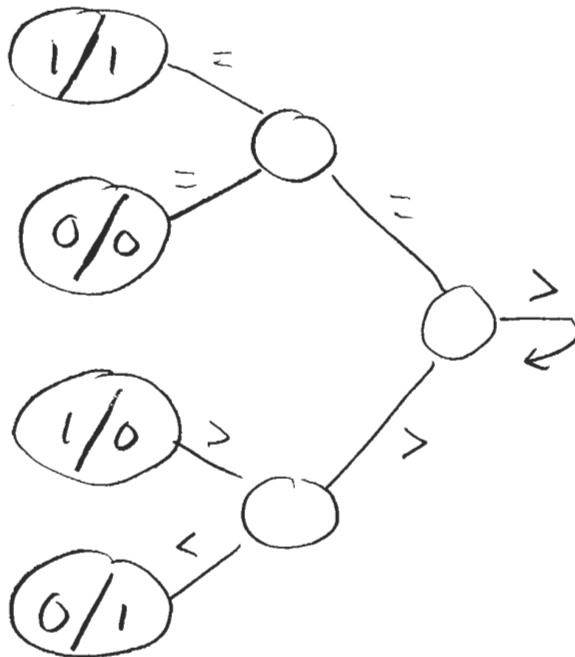
6.896  
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## Comparing 2 k-bit words



$\Theta(k)$  steps to compute  
Sort in  $\Theta(Nk)$  time

## Faster comparison - binary tree

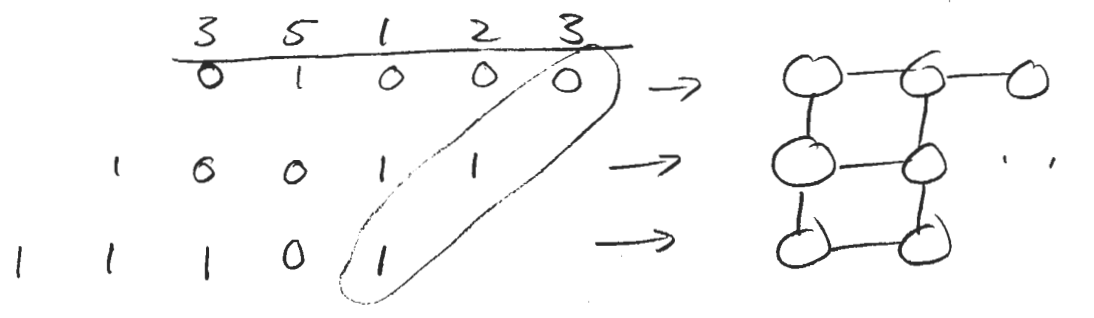


$\Theta(\lg k)$  compare  
Sort in  $\Theta(N \lg k)$   
 $\lg = \log_2$

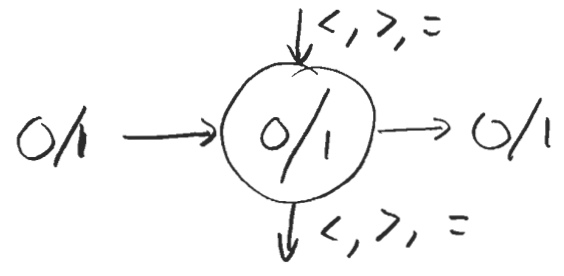
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2/4/04  
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# Pipelining

- compare while sorting
- stagger bits of input.



Each processor:



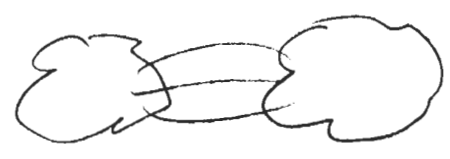
Time =  $\Theta(N+k)$  bit steps.  
Can we do better on  $N \times k$  array?

## Lower bounds

1. I/O bandwidth.  
 $Nk$  bits to input at  $k$  places  $\Rightarrow \Omega(N)$  steps.
2. Network diameter  
 $\Omega(N+k)$
3. Communication bandwidth (bisection width).

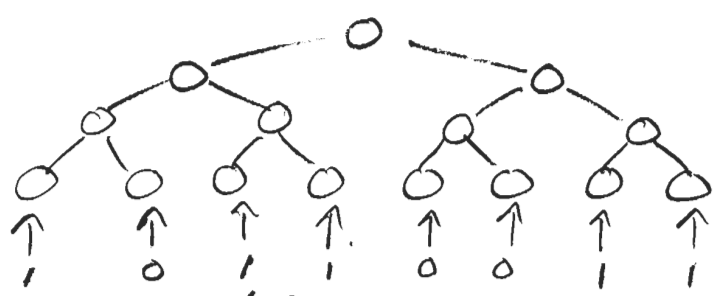
$$T \geq \frac{\# \text{ bits crossing cut}}{\text{size of cut}}$$

$$T \geq \frac{\Theta(Nk)}{\Theta(k)} = \Theta(N).$$



6.896  
2/4/04  
L 1.5

Problem.  $N$  1-bit #'s input at  $N$  leaves of complete binary tree. Time to sort?

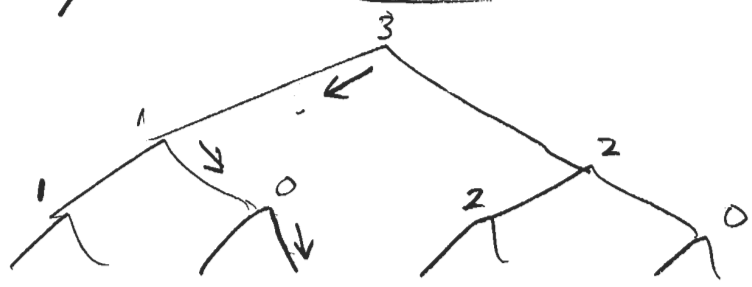


I/O :  $T \geq N / N = 1$   
 Diam :  $T \geq 2 \lg N$   
 BW :  $T \geq \Theta(N) / 1 = \Theta(N)$

$\therefore$  Sorting  $N$  1-bit #'s takes  $\Omega(N)$  time on CBT.

Wrong! Can sort in  $O(\lg N)$  time!

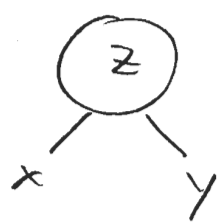
Idea: Only need to count # 0's.



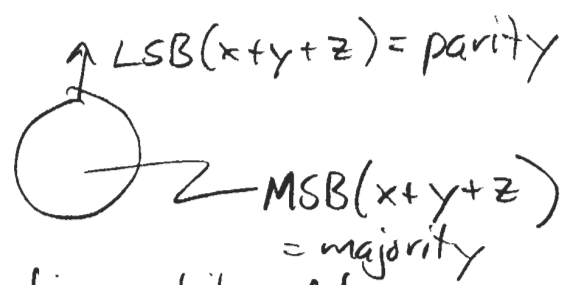
Input 1 0 1 1 0 0 1 1  
 Output 0 0 0 1 1 1 1 1

Sum 0's upward, select downward.

1-bit summer



$\Rightarrow$



Pipeline in bit model.

Q. Why doesn't BW lower bound hold?  
 A. Didn't really need to ship  $\Theta(N)$  bits across bisection.  
 Could encode info more compactly.  
 Careful!