MASSACHUSETTS INSTITUTE OF TECHNOLOGY SLOAN SCHOOL OF MANAGEMENT

# **<u>15.565</u>** Integrating Information Systems:

Technology, Strategy, and Organizational Factors

# **<u>15.578</u>** Global Information Systems:

**Communications & Connectivity Among Information Systems** 

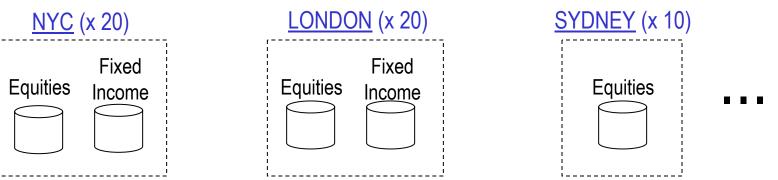
Spring 2002

Lecture 15

## DISTRIBUTED HETEROGENEOUS DATABASES

### **NEXT GENERATION GLOBAL FOREIGN EXCHANGE SYSTEM (FXX)**

#### TRADING SYSTEMS

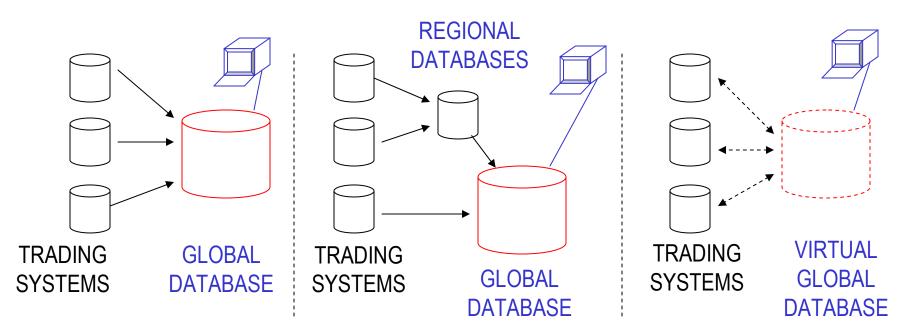


- EXAMPLE QUESTIONS:
  - WHAT ARE OUR TOTAL HOLDINGS OF "IBM"?
  - WHAT ARE TOTAL ASSETS OWNED BY & EXPOSURES OF "LONG-TERM CAPITAL"?
  - WHAT ARE TOTAL ASSETS & EXPOSURES ASSOCIATED WITH "BRAZIL" ECONOMY?

#### • OTHER SIMILAR SITUATIONS:

- CUSTOMER RELATIONSHIPS
- PROFITABILITY ANALYSIS (BY PRODUCT, DIVISION, ETC.)
- ASSET VISIBILITY (INVENTORY, IN-TRANSIT, ETC.)
- DATA WAREHOUSES (IN GENERAL)
- WHY NOT <u>ONE</u> <u>BIG</u> DATABASE?

### **POSSIBLE "GLOBAL DATABASE" CONFIGURATIONS**

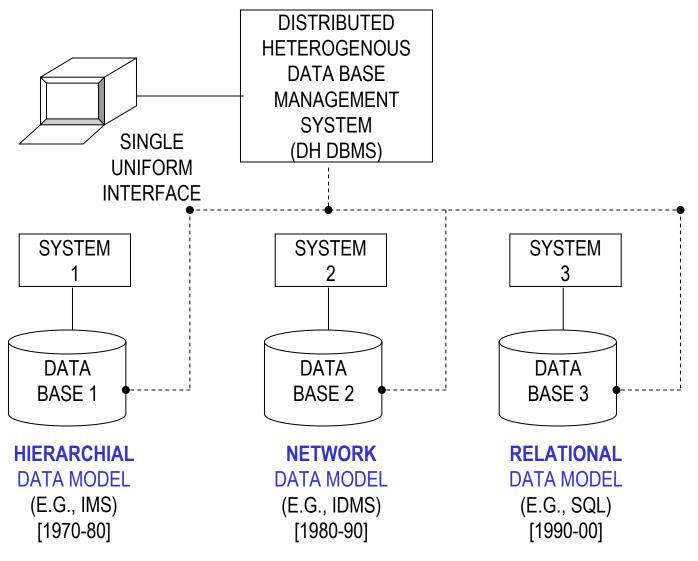


- ADVANTAGES AND DISADVANTAGES OF EACH APPROACH
- FURTHER COMPLICATIONS
  - THE TRADING SYSTEMS USE DIFFERING HARDWARE AND OPERATING SYSTEMS (MAINFRAMES, MINI, PC).
  - THE TRADING SYSTEMS USE DIFFERING DATABASE MANAGEMENT SYSTEM (DBMS) SOFTWARE

#### $\rightarrow \text{HETEROGENEOUS ENVIRONMENT}$

#### **DISTRIBUTED HETEROGENOUS DATA BASE MANAGEMENT SYSTEMS**

#### (VIRTUAL GLOBAL DATABASE)

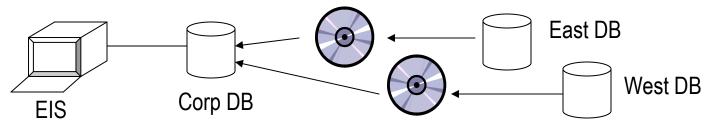


• "LEGACY" SYSTEMS

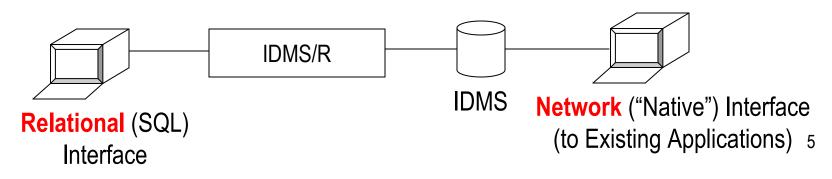
## **VARIOUS APPROACHES TO DISTRIBUTED DATABASES**

1. BATCH MODE -- PERIODIC EXTRACTS, NO AD HOC QUERY

--DOWN-LOAD, SHADOW DATABASES, SNEAKER-NET



- 2. HETEROGENEOUS INTERFACE -- ONLINE ACCESS TO SINGLE DATABASE
  - -- NO DATABASE INTEGRATION
  - -- ACCESS USES ONE DATA MODEL; DBMS ANOTHER
  - -- E.G., RELATIONAL INTERFACE TO NETWORK SYSTEM (IDMS/R)



## **VARIOUS APPROACHES TO DISTRIBUTED**

### **DATABASES** (continued)

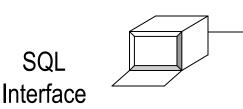
**INTEROPERABILITY** -- ACCESS TO MULTIPLE DATABASES; 3.

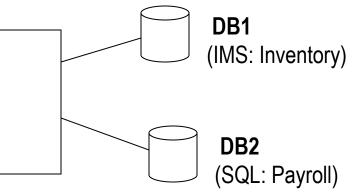
### SIMILAR OR DIFFERENT DBMS'S

-- NO LOCATION TRANSPARENCY

SQL

-- NON-INTEGRATED SCHEMAS

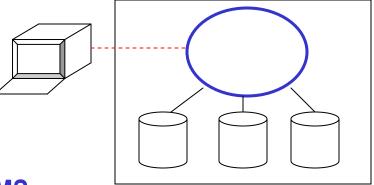




4. GATEWAY -- PARTIAL INTEGRATION: MOSTLY HOMOGENEOUS DBMS'S -- GATEWAY TO HETEROGENEOUS DBMS'S IBM -- NO REAL DISTRIBUTED QUERIES DB2 Oracle/IBM SQL Oracle/HP 6 Interface Oracle\*

## VARIOUS APPROACHES TO DISTRIBUTED DATABASES (continued)

- 5. HETEROGENEOUS DDBMS
  - -- SINGLE GLOBAL SCHEMA
  - -- COMMON DATA MODEL AND QUERY LANGUAGE
  - -- INTEGRATED SCHEMA
  - -- TRUE DISTRIBUTED QUERY
  - -- E.G., MULTIBASE



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- 6. FEDERATED HETEROGENEOUS DDBMS
  - -- <u>MULTIPLE</u> SCHEMAS
  - -- AUTONOMY OF COMPONENT DATABASES
  - -- TRUE DISTRIBUTED QUERY

#### TRUE HETEROGENEOUS DATABASES = # 5 AND # 6

## **TASKS PERFORMED BY DHDBMS**

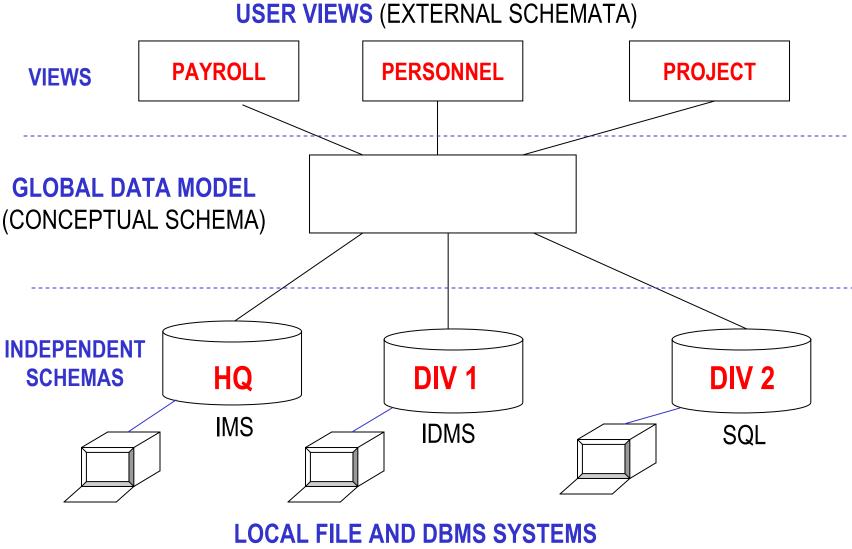
- 1. ACCEPT STANDARD QUERY AND UPDATE REQUESTS
- 2. AUTOMATICALLY DIVIDE UP QUERY FOR INDIVIDUAL DATABASES
- 3. TRANSFORM INTO SUBQUERIES IN APPROPRIATE LOCAL DBMS LANGUAGE
- 4. FORMULATE SEQUENCE OF EXECUTION AND DATA MOVEMENT (Optimization)
- 5. RESOLVE INCOMPATIBILITIES BETWEEN DATABASES (E.G., DOLLARS vs. POUNDS vs. YEN)
- 6. RESOLVE INCONSISTENCIES IN COPIES OF SAME INFORMATION (E.G., 77 MASS AVE. vs. 50 MEMORIAL DRIVE)
- 7. COMBINE RETRIEVED DATA INTO CONSOLIDATED RESPONSE

## **COMMENTS AND CONSTRAINTS ON A DHDBMS**

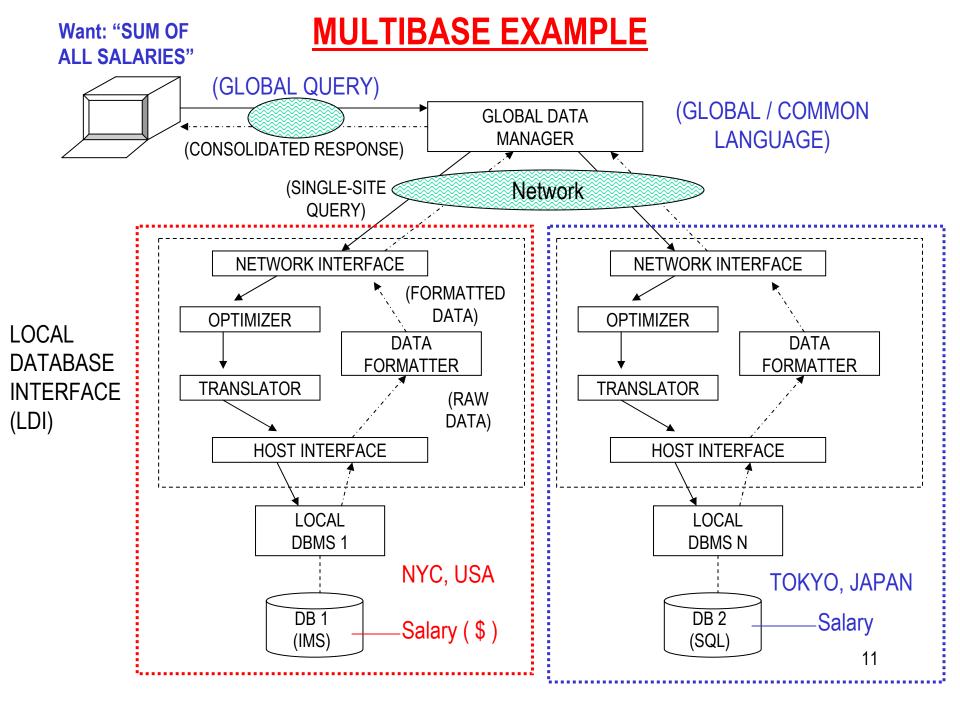
- 1. COMPONENTS ARE <u>EXISTING</u> DBMS SYSTEMS, NOT DESIGNED FOR DDBMS PURPOSES
- 2. COMPONENTS CANNOT BE EASILY MODIFIED– WHY?

- 3. MUST BE ABLE TO ADD AND DELETE COMPONENTS TO DHDBMS
- 4. NOTION OF "NON-INTRUSIVE"

### **THREE SCHEMA APPROACH TO DHDBMS**



(INTERNAL SCHEMATA)



### **MULTIBASE CAPABILITIES**

- MULTIBASE:
  - UNIFORM, RETRIEVAL INTERFACE (DAPLEX)
  - SINGLE QUERY LANGUAGE
  - PRE-EXISTING, HETEROGENEOUS DISTRIBUTED DATABASES
- OBJECTIVES: GENERALITY, COMPATIBILITY, AND EXTENSIBILITY
- NON-INTRUSIVE: NO CHANGES NEEDED TO UNDERLYING DBMS'S, FILE SYSTEMS, OR APPLICATIONS
- CAPABILITIES:
  - SCHEMA INTEGRATION -- DEFINE GLOBAL VIEW FROM LOCAL SCHEMAS
  - DATA INCOMPATIBILITY HANDLING -- MAINTAIN AND ENFORCE RULES FOR RECONCILING DIFFERENCES
  - QUERY OPTIMIZATION -- OPTIMIZE LOCALLY AND GLOBALLY (COMPENSATING FOR LOCAL SYSTEM LIMITATIONS)
  - QUERY TRANSLATION -- TRANSLATE FROM DAPLEX TO LOCAL QUERY LANGUAGES

## **SCHEMA INCOMPATIBILITIES**

- PROBLEMS (TO BE DISCUSSED IN NEXT SESSION)
  - -- DIFFERENCES IN:
    - -- NAMING CONVENTIONS
    - -- UNDERLYING DATA STRUCTURES
    - -- REPRESENTATIONS (Male/Female, Man/Woman, M/F, 0/1)
    - -- SCALES / UNITS
  - -- MISSING DATA
  - -- CONFLICTING DATA VALUES
- SOLUTIONS: (MULTIBASE)
  - A VIEW MECHANISM, WITH RENAMING, LOGICAL RESTRUCTURING, AND SCALE CONVERSION, CAN SOLVE "DIFFERENCE" PROBLEMS
  - AUXILIARY DATABASE CAN BE ATTACHED TO THE UNIFIED SCHEMA TO SOLVE OTHER TWO TYPES OF PROBLEMS

### **QUERY PROCESSING**

- PROBLEMS
  - -- USE PARALLEL PROCESSING & MINIMIZE PROCESSING AND COMMUNICATIONS COSTS
  - -- ISSUES:
    - -- LOCAL QUERY COSTS DIFFER AT DIFFERENT SITES
    - -- MANY DBMS'S DON'T SUPPORT CREATION OF TEMPORARY DATABASES BY GDM
    - -- LOCAL QUERYING CAPABILITIES MAY BE LIMITED
    - -- LOCAL DBMS'S HAVE VARYING OPTIMIZATION CAPABILITIES
- SOLUTIONS / ALTERNATIVES
  - -- <u>ALTERNATIVE 1</u>: GLOBAL QUERY --> <u>SMALLEST</u> POSSIBLE SINGLE SITE SUBQUERIES
    - -- COULD BE MULTIPLE SUBQUERIES AT A SITE, MERGE RESULTS AT GDM
    - -- MORE GDM WORK AND COMMUNICATION, BUT SIMPLE
  - -- <u>ALTERNATIVE 2</u>: GLOBAL QUERY --> <u>LARGEST</u> POSSIBLE SINGLE SITE SUBQUERIES
    - -- LESS MERGING WORK, LESS COMMUNICATION
    - -- MORE LDI SOPHISTICATION NEEDED
  - -- <u>ALTERNATIVE 3</u>: GENERATE EFFICIENT PLAN, LDI'S PARTICIPATE IN GLOBAL OPTIMIZE
    - -- SEND PARTIAL RESULTS TO GDM OR OTHER LDIS
    - -- LDI'S NEED TO SUPPORT SORTING, DUPLICATE REMOVAL, MERGING TEMPORARIES
    - -- LESS WORK BY GDM, LESS COMMUNICATION
    - -- COMPLEX GDM AND LDI'S

## TRANSACTION MANAGEMENT (COMPLEX UPDATES)

- PROBLEMS
  - -- DIFFERENT DBMS'S -- DIFFERENT CONCURRENCY CONTROL METHODS
  - -- IMPLEMENTATION DIFFERENCES:
    - -- DEADLOCK HANDLING & LOCKING GRANULARITY
  - -- COMMIT PROTOCOLS MAY VARY
  - -- GLOBAL DEADLOCK DETECTION DIFFICULT:
    - -- LOCAL PROCESSES DON'T KNOW ABOUT NON-LOCAL TRANSACTIONS
    - -- GLOBAL PROCESSES DON'T KNOW ABOUT LOCAL TRANSACTIONS
- SOLUTIONS/ALTERNATIVES
  - -- ALLOW READS ONLY--NO WRITES. NO CONCURRENCY CONTROL PROBLEM.
  - -- ALLOW SINGLE SITE UPDATES ONLY
    - -- USE LOCAL CONCURRENCY CONTROL MECHANISMS
  - -- ASSUME ALL DBMS'S PROVIDE LOCKING AND TWO-PHASE COMMIT.
    - -- DETECT CONFLICTS BY TIME OUT.
  - -- ALLOW MULTI-SITE UPDATES.
    - -- MAY REQUIRE MODIFICATIONS TO LOCAL DBMS'S AND OS'S. <sup>15</sup>

### **SUMMARY**

- DISTRIBUTED HETEROGENEOUS DATA BASES ARE A REALITY IN MOST ORGANIZATIONS -- FOR BETTER OR WORSE.
- MOST APPROACHES TO HANDLE AGGREGATION AND COORDINATED ACCESS -- HAVE BEEN <u>AD-HOC</u>.
- DISTRIBUTED HETEROGENEOUS DATABASE MANAGEMENT SYSTEMS (DHDBMS) ARE EMERGING TO ADDRESS THESE NEEDS.